

Recall

①

- $H(X) = -\sum P(X=x) \log_2 P(X=x)$
 - $H(X|Y) = H(X, Y) - H(Y)$
 $= \mathbb{E}_y H(X|Y=y)$
 - $0 \leq H(X) \leq \log_2 |\text{supp}(X)|$ " $=$ " iff X unif
 - $H(X, Y) \leq H(X) + H(Y)$
 - $H(X|Y) \leq H(X)$
 - $H(X, f(X, Y) | Y) = H(X|Y) = H(X|Y, f(Y))$ for deterministic f
 - $H(f(X) | Y) \leq H(X|Y) \leq H(X|f(Y))$
 - Bit-wise entropy calculation
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Graph hom

homomorphism

$$\text{Hom}(H, G) = \{ f: V(H) \rightarrow V(G) \mid f(u)f(v) \in E(G) \forall uv \in E(H) \}$$

$$\text{hom}(H, G) = |\text{Hom}(H, G)|$$

hom density

$$t(H, G) = \frac{t(H, G)}{|V(G)|^{|V(H)|}}$$

Sidorenko

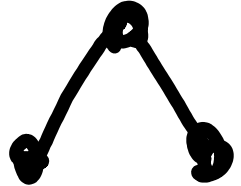
Def A graph H is Sidorenko if

$$t(H, G) \geq t(K_2, G)^{|E(H)|} \quad \forall G$$

Remark • If $G = G(n, p)$, then $\mathbb{E} t(H, G) = p^{|E(H)|}$
 $= (\mathbb{E} t(e, G))^{|E(H)|}$

Roughly speaking, Sidorenko property says random graph minimizes the density of H

- If H NOT bipartite, then H NOT Sidorenko
 pf take $G = e$ $\square$

Thm P_2 :  side.

$$\text{pf } \text{hom}(P_2, G) = \sum_v \deg(v)^2 \geq \frac{(\sum \deg(v))^2}{n}, \quad n = |V(G)|$$

$$= \frac{(2m)^2}{n}, \quad m = |E(G)|$$

$$t(K_2, G) \leq \frac{(2m)^2}{n} / n^3 = \left(\frac{2m}{n^2}\right)^2 = t(e, G)^2$$

Alt pf Take $XY \in E(G)$ uar.

Condition on Y , resample X indep. and get Z

$$(i.e. P(Z=z, X=x | Y=y) = P(X=z | Y=y) \cdot P(X=x | Y=y))$$

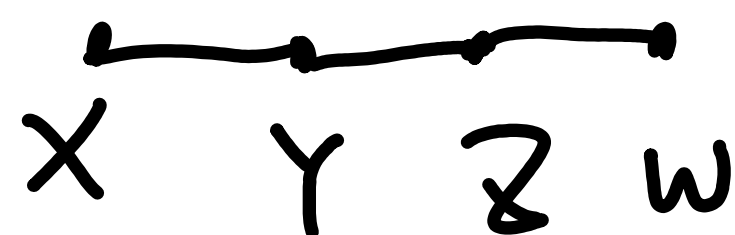
$$\Rightarrow H(X, Z | Y) = H(X | Y) + H(Z | Y) = 2H(X | Y)$$

$$\underbrace{H(XYZ)}_{\substack{\uparrow \\ \log \text{hom}(P_2, G)}} = 2H(X | Y) + H(Y) = \underbrace{2H(XY)}_{\log \text{hom}(K_2, G)} - \underbrace{H(Y)}_{\log_2 n} \quad (n = |V(G)|)$$

Thm P_3 side

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pf Take $XY \in E(G)$ u.a.r.



Condition on Y , resample X indep. get Z

Condition on Z , resample Y indep. get W

$$H(X, Y, Z, W) \leq \log_2(\text{hom}(P_3, G))$$

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$$\frac{H(X, Y, Z) + H(W | X, Y, Z)}{H(X, Y) + H(Z | Y) + H(W | Z)}$$

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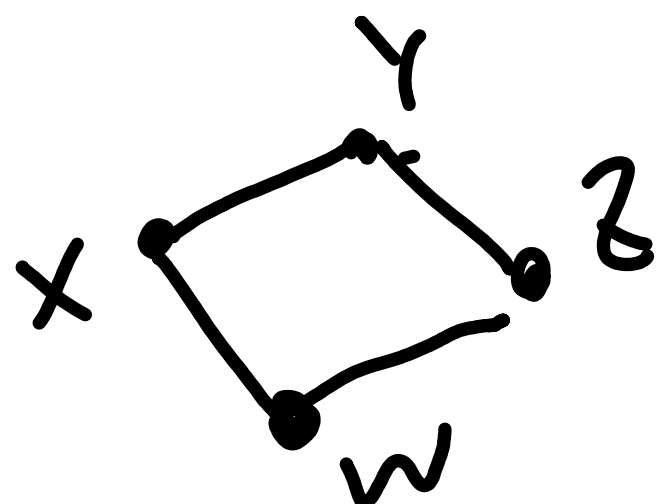
$$= \frac{H(X, Y) + H(Y, Z) + H(Z, W)}{H(X, Y) + H(Z | Y) + H(W | Z)} = \frac{H(X, Y) + H(Y, Z) + H(Z, W)}{H(X, Y) + H(Z | Y) + H(W | Z)}$$

$$\frac{-H(Y) - H(Z)}{-2 \log_2 n} \parallel 3 \log_2(2m)$$

$$\Rightarrow \frac{\text{hom}(P_3, G)}{n^4} \geq \frac{(2m)^3}{n^6} \quad \square$$

Thm C_4 side.

pf



Condition on X, Z , resample Y indep. get W

$$H(X, Y, Z, W) \leq \log_2(\text{hom}(P_4, G))$$

"

$$H(X, Y, Z) + H(W | X, Z) = 2H(X, Y, Z) - H(X, Z)$$

$$= 4H(X, Y) - 2H(Y) - H(X, Z)$$

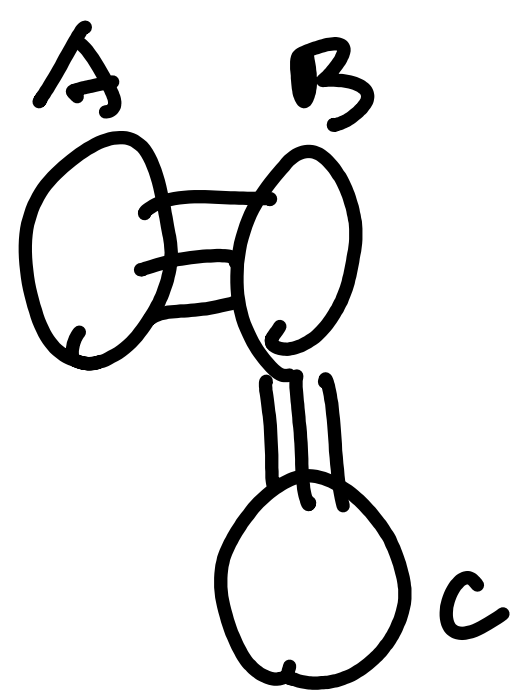
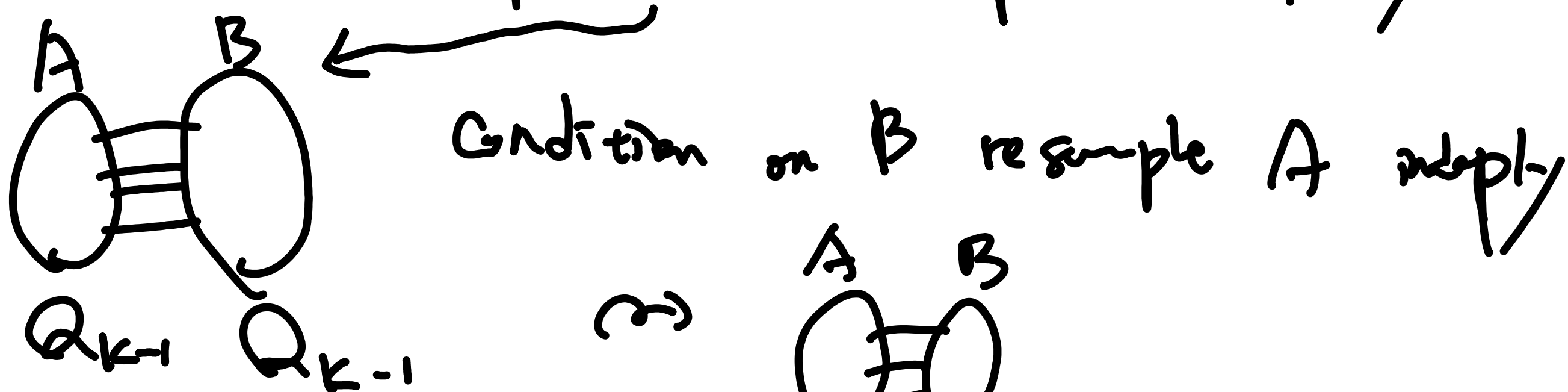
$$H(X, Z) \leq H(X) + H(Z) \leq 2 \log_2 n$$

$$\Rightarrow \frac{\text{hom}(C_4, G)}{n^4} \geq \frac{(2m)^4}{n^8} \quad \square$$

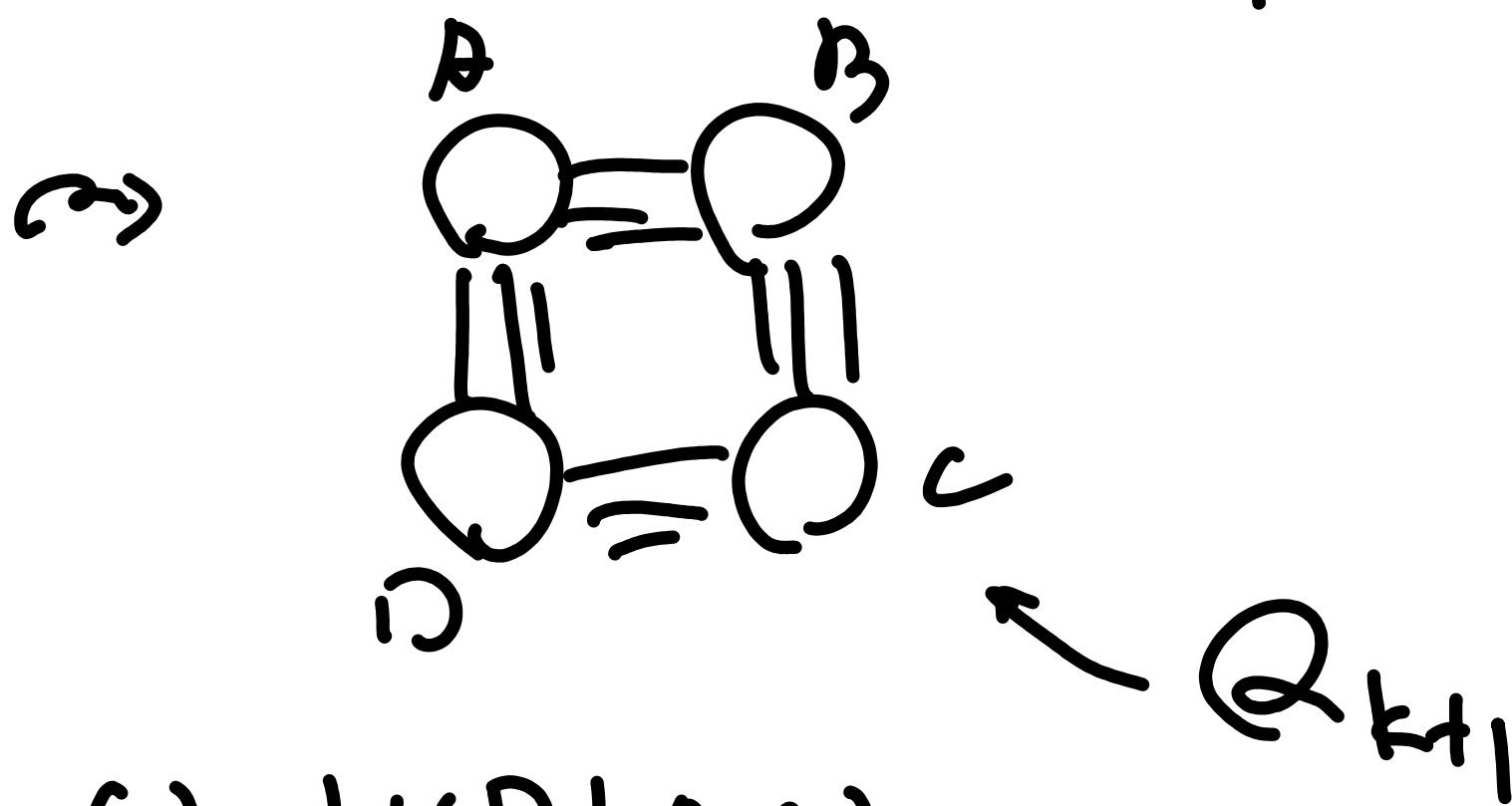
Thm $Q_0 = \bullet$, $Q_1 = \text{---}$, $Q_2 = \square$, $Q_3 = \text{cube}$

Q_n side.

Pf Assume we sampled Q_k so that the marginal distribution of $A \sim B$. Sample Q_{k+1} by



Condition on A, C , resample B indep.



$$\Rightarrow H(Q_{k+1}) = H(A, B, C) + H(D | A, C)$$

$$\geq 4H(A, B) - 4H(A) = 4H(Q_k) - 4H(Q_{k-1})$$

$$\text{Let } D_k = H(Q_k) - 2^k \log_2 n \leq \log_2 \text{hom}(Q_k, G)$$

$$\Rightarrow D_{k+1} \geq 4D_k - 4D_{k-1}$$

$$D_n \geq 4D_{n-1} - 4D_{n-2}$$

$$4D_{n-1} \geq 16D_{n-2} - 16D_{n-3}$$

$\vdots$

$$(k+1) \cdot 2^k D_{n-k} \geq (k+1) 2^{k+2} D_{n-k-1} - (k+1) 2^{k+2} D_{n-k-2}$$

$\vdots$

+)

$$D_n \geq (n-1) 2^n D_1 - (n-1) 2^n D_0$$

$$= (n-1) 2^n \log_2 \left(\frac{2m}{n} \right)$$

$$D_1 = \log_2 \left(\frac{2m}{n} \right), D_0 = 0$$

hom density of trees

Thm (Sidorenko)

Let T be a tree w/ k vtxs,

$$t(T, G) \leq t(K_{1,k-1}, G) \quad \forall G$$

$$\text{hom}(T, G) \leq \text{hom}(K_{1,k-1}, G)$$

Pf Claim Fix T, G , there exists T' : tree w/ k vtxs

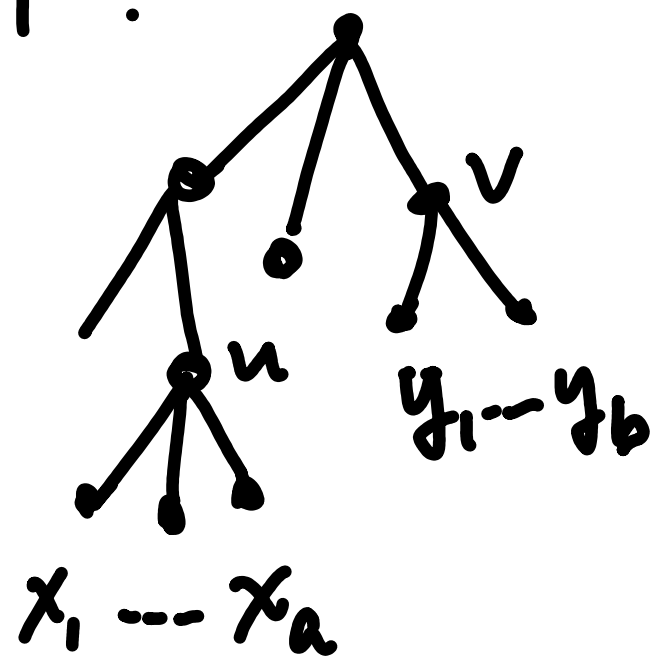
and has more leaves than T s.t.

$$\text{hom}(T, G) \leq \text{hom}(T', G)$$

Pf Pick $X_T = (X_v)_{v \in V(T)} \in \text{Hom}(T, G)$ unif at random.

Pick two vtxs u, v that are distance one from some leaves ^(b)
 $\text{in } T$

T :



$$H(X_T) = H(X_T \setminus \{x_1, \dots, x_a, y_1, \dots, y_b\}) + H(X_{x_1}, \dots, X_{x_a}, X_{y_1}, \dots, X_{y_b} | X_{\tilde{T}})$$

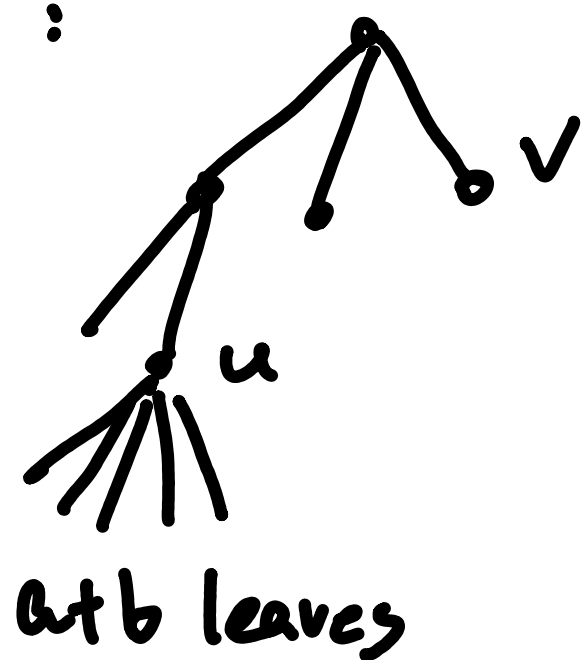
(actually $\Rightarrow$)

$$\leq H(X_{\tilde{T}}) + a H(X_x | X_{\tilde{T}}) + b H(X_y | X_{\tilde{T}})$$

$$\leq H(X_{\tilde{T}}) + a H(X_x | X_u) + b H(X_y | X_v)$$

Case 1 $H(X_x | X_u) \geq H(X_y | X_v)$

T' :



$$H(X_T) \leq H(X_{\tilde{T}}) + (a+b) H(X_x | X_u)$$

resample $X_{x_{a+i}} | X_u \sim X_{x_i} | X_u$, $i=1, \dots, b$
 c.i.

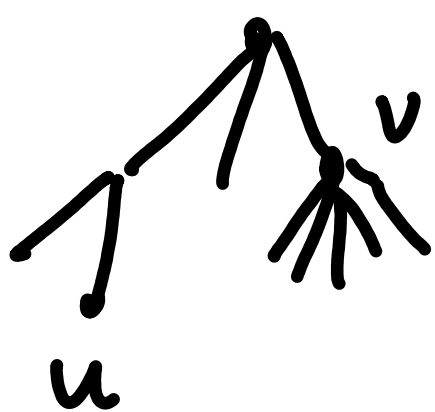
$$H(\tilde{T}, X_{x_1}, \dots, X_{x_{a+b}}) = H(X_{\tilde{T}}) + (a+b) H(X_x | X_u)$$

↑
Supported on $\text{Hom}(T', G)$

$$\Rightarrow \log_2 \text{hom}(T, G) = H(X_T) \leq H(\tilde{T}) \leq \log_2 \text{Hom}(T', G)$$

Case 2 $H(X_x | X_u) \leq H(X_y | X_v)$

T' :



□

Shearer's inequality

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Thm $H(X, Y, Z) \leq H(X, Y) + H(Y, Z) + H(X, Z)$

pf $H(X, Y, Z) = H(X, Y) + \underline{H(Z | X, Y)}$

$$= H(Y, Z) + H(X | Y, Z)$$

$$= H(X, Z) + \underline{H(Y | X, Z)}$$

$$H(Z | XY) + H(X | YZ) + H(Y | XZ) \leq H(Z) + H(X | Z) + H(Y | XZ) \\ = H(XYZ)$$

Thm Suppose $A_1, \dots, A_m \subseteq [n]$ covers i at least k times $\forall i \in [n]$
then $kH(X_1, \dots, X_n) \leq \sum_{i=1}^m H(X_{A_i})$

where $H(X_A) = (X_{a_1}, \dots, X_{a_t})$ if $A = \{a_1, \dots, a_t\}$

pf Submodularity $\Rightarrow H(X_{A_i}) + H(X_{A_j}) \geq H(X_{A_i \cup A_j}) + H(X_{A_i \cap A_j})$

Replace A_i, A_j by $A_i \cup A_j, A_i \cap A_j$, RHS decreases
the family still covers i $\geq k$ times $\forall i \in [n]$

Keep applying this move until we can't (stop in finite time)
why?

WLOG we get $A_1 \supseteq \dots \supseteq A_k \supseteq \dots \supseteq A_m$

Since it covers i $\geq k$ times, $\bigcap_{i=1}^k A_i = \dots = A_k = [n]$

$$\Rightarrow \text{RHS} \geq kH(X_{[n]})$$