

In the following examples, we for simplicity use the approximation  $\arcsin(x) \approx x$ . In reality, one should use proper (inequality) estimates, often using different estimates for  $x < 0$  &  $x > 0$ . There are plenty of interesting applications where this is enough. There are also a handful of applications where one has to both at higher order terms ( $O(x^3)$ -terms).

Our first example is proving  $sp(G) = \Omega(m^{4/5})$  whenever  $G$  is triangle-free. First, where is  $4/5$  coming from? A good guess is that pseudorandom triangle-free graphs would be good choice. That is, we have  $|E|, \Delta_2 \sim \sqrt{d}$ .

$G$  triangle-free  $\Rightarrow 0 = \text{tr}(A^3) \geq d^3 + n \Delta_2^3 \Rightarrow d^3 \sim n d^{3/2}$   
 $\Rightarrow$  need  $d \leq n^{2/3}$  for pseudorandom  $\Delta$ -free graphs to exist.

For pseudorandom graphs, we'd also have  $sp(G) \sim n \sqrt{d} \sim n^{4/5}$ .

and  $m = \frac{1}{2} n d \sim n^{5/3} \approx \frac{2}{3} \frac{m^{4/5}}{n^{1/5}} \Rightarrow sp(G) \sim m^{4/5}$ .

In fact, Alon proved that pseudorandom  $\Delta$ -free graphs exist with  $d \sim n^{2/3}$  for infinitely many  $n$ .

Theorem 29 (Alon)  $\exists$  infinite sequence  $q_k \rightarrow \infty$  for which there exists  $G_k$  with  $n_k \sim q_k^3$ ,  $d_k \sim q_k^2$ ,  $\Delta_2(G_k) \sim q_k$  which are  $\Delta$ -free (Thus  $sp(G_k) \sim q_k^4 \sim m_k^{4/5}$ ,  $m_k = \frac{1}{2} n_k d_k \sim q_k^5$ ).

Lemma 30 Let  $G$  be  $\Delta$ -free graph with degree sequence  $d_1 \rightarrow d_n$ . Then  $sp(G) = \Omega\left(\sum_{i=1}^n \sqrt{d_i}\right)$ .

Think of  $G$   $d$ -regular  $\Rightarrow sp(G) = \Omega(n \sqrt{d})$ .

Proof: Choose  $x_u \in \mathbb{R}^{V(G)}$  with

$$x_u(v) = \begin{cases} 1 & \text{if } u=v \\ 0 & \text{if } uv \notin E \text{ \& } u \neq v \\ -1/\sqrt{d_u} & \text{if } uv \in E. \end{cases}$$

Given  $u, v \in E$ , we have  $x_u(s)x_v(s) \neq 0 \Rightarrow s \in N[u] \cap N[v]$ .

( $N[w] = \{w\} \cup N(w)$ ). Now,  $N(u) \cap N(v) = \emptyset$  as  $G$  is  $\Delta$ -free.

$\Rightarrow x_u(s)x_v(s) \neq 0 \Rightarrow s \in \{u, v\}$ .

Thus  $\langle x_u, x_v \rangle = -1/\sqrt{d_u} - 1/\sqrt{d_v}$ . Also,  $\|x_u\|_2^2 = 2 \forall u$ .

Thus  $sp(G) \geq -\frac{1}{\Omega} \sum_{u,v \in E} \arcsin\left(\frac{\langle x_u, x_v \rangle}{\|x_u\|_2 \|x_v\|_2}\right)$ .

$$= \Omega \left( \sum_{u,v \in E} \left( \frac{1}{\sqrt{d_u}} + \frac{1}{\sqrt{d_v}} \right) \right) = \Omega \left( \sum_u \sqrt{d_u} \right) \square.$$

This bound allows us to deal with graphs with "not too many large degrees" (in fact with graphs that are  $d$ -degenerate for  $d \sim m^{2/5}$ ). But we need other argument for graphs that are "too dense" (i.e. ~~are~~ more precisely, for those graphs that contain subgraph with large min degree).

Lemma 31: Suppose  $G$  has  $m$  edges and is  $d$ -degenerate for  $d = m^{2/5}$ . Then  $sp(G) = \Omega(m^{4/5})$ .

Proof: Pick  $v_1 \rightarrow \dots \rightarrow v_n$  so that  $d_{<i} := \left\{ j : v_i v_j \in E \text{ \& } v_j < i \right\} < d \forall i$ .

Thus,  $sp(G) = \Omega \left( \sum_{i=2}^n \sqrt{d_{<i}} \right) \geq \Omega \left( \sum_{i=1}^n \sqrt{d_{<i}} \right) \geq \frac{\Omega \left( \sum_{i=1}^n d_{<i} \right)}{\sqrt{d}}$ .

$$= \frac{\Omega(m)}{\sqrt{d}} = \Omega(m^{4/5}). \text{ as } \sum_{i=1}^n d_{<i} = m. \square$$

Exercise 32 Suppose  $H_1, \dots, H_t$  are  $r$  disjoint <sup>induced</sup> subgraphs of  $G$ .  
 Then  $sp(G) \geq \sum_{i=1}^t sp(H_i)$ . In particular,  $sp(G) \geq sp(H_1)$ .

Lemma 33: Suppose that  $G$  is not  $d$ -degenerate for  $d = m^{2/5}$  and  $e(G) = m$ . Then  $sp(G) = \Omega(m^{4/5})$ .

Proof: Let  $H'$  be induced subgraph with  $\delta(H') \geq d$ .

Let  $R \subseteq V(H')$  be a random subset obtained by choosing  $v_1, \dots, v_r$  from  $H'$  uniformly at random (allowing repeats), where  $r = \lfloor 2|H'|/d \rfloor$ .

Let  $H''$  denote the induced subgraph of  ~~$H'$~~   $H'$  so that  $v \in H'' \iff \forall i \in [r] \ vv_i \in E$ .

Then  $P(v \notin H'') \leq \left(1 - \frac{d}{|H'|}\right)^r \leq \exp\left(-\frac{dr}{|H'|}\right) \leq \exp(-2) < 1/4$ .

$\Rightarrow P(\forall v \in E(H'')) > 1 - 1/4 - 1/4 = 1/2$  for  $\forall v \in E(H)$ .

$\Rightarrow \mathbb{E}[e(H'')] > 1/2 e(H') \geq 1/4 d |H'|$ .

We also claim that  $\chi(H'') \leq r$ . Indeed, let  $\chi: V(H'') \rightarrow [r]$  be defined as  $\chi(u) = \min\{i: uv_i \in E\}$ .

Suppose that  $\chi(u) = \chi(v) = i$  &  $uv \in E(H'')$ .

Then  $uv, uv_i, vv_i \in E(G) \Rightarrow u, v, v_i$  form  $\Delta$  in  $G$ .

Thus  $\chi(H'') \leq r \Rightarrow sp(H'') \geq \Omega\left(\frac{e(H'')}{\chi(H'')}\right)$ .

$= \Omega\left(\frac{d|H'|}{|H'|/d}\right) = \Omega(d^2) = \Omega(m^{4/5})$ .

Thus  $sp(G) = \Omega(m^{4/5})$  follows from Exercise 32.  $\square$

Theorem 34 (Alon) Suppose that  $G$  is  $\Delta$ -free.  
Then  $sp(G) = \Omega(m^{4/5})$ , and the exponent  $4/5$  is tight.

Proof: Tightness follows from Theorem 29. If  $G$  is  $m^{2/5}$ -degenerate, Lemma 31 implies the result. Otherwise, Lemma 33 implies the result.

GW for disc<sup>+</sup>

Suppose that for each  $u \in V(G)$   $x_u$  is a unit vector. Then

$$\text{disc}^+(G) \geq \frac{1}{2\gamma} \left( \sum_{u \in E} \left( \frac{\gamma}{2} + \alpha \cos(\langle x_u, x_v \rangle) \right) \right) - \frac{p}{2} \sum_{u \neq v} \left( \frac{\gamma}{2} + \alpha \cos(\langle x_u, x_v \rangle) \right).$$

Note that the constant terms simply cancel out.

Again, we use the approximation

$$\text{disc}^+(G) = \Omega \left( \sum_{u \in E} \langle x_u, x_v \rangle \right) - O \left( p \cdot \sum_{\substack{u \neq v \\ u, v \in E}} \langle x_u, x_v \rangle \right).$$

(Here, we are assuming that  $p \leq 1 - \epsilon$ .)

Theorem 36 (R., Sudakov, Tomon '20)

$$\text{disc}^+(G) \leq \begin{cases} \Omega(\sqrt{d} \cdot n) & \text{when } 1 \leq d \leq n^{2/3} \\ \Omega(n^2/d) & \text{when } n^{2/3} \leq d \leq \frac{1}{2}n. \end{cases}$$

We only prove Thm 36 for regular graphs (Actually, the proof works if  $\Delta \leq C \cdot d$ ,  $C$  constant)

Proof: Choose  $x_u(v) = \begin{cases} 1 & \text{if } u=v \\ \alpha & \text{if } u \neq v \in E, \text{ or } \alpha \leq 1/\sqrt{d} \text{ to be determined later} \\ 0 & \text{otherwise.} \end{cases}$

Then  $\|x_u\|_2^2 = 1 + d\alpha^2 \leq 2$  as  $\alpha \leq 1/\sqrt{d}$ .

If  $uv \in E$ , then  $\langle x_u, x_v \rangle = 2\alpha + \alpha^2(N(u) \cap N(v)) \geq 2\alpha$ .  
 Otherwise,  $\langle x_u, x_v \rangle = \alpha^2(N(u) \cap N(v))$ .

$$\text{Thus } \sum_{uv \in E} \langle x_u, x_v \rangle \leq \sum_{u \neq v} \alpha^2 |N(u) \cap N(v)| \leq \alpha^2 nd(d-1)$$

as each  $w \in V$  contributes ~~at most~~ to exactly  $d(d-1)$  terms in the sum (for any  $(u, v) \in N(w) \times N(w)$  with  $u \neq v$ )

$$\begin{aligned} \text{Thus, } \text{disc}^1(G) &= \Omega(\alpha nd) - O(p \cdot nd^2 \alpha^2) \\ &= \Omega(\alpha nd) - O(d^3 \alpha^2). \end{aligned}$$

(i) If  $d \leq n^{2/3}$ , choose  $\alpha = c\sqrt{d}$  (for sufficiently small constant).

Then  $\alpha nd = c\sqrt{d}n$  and  $d^3 \alpha^2 = c^2 d^2 = c^2 \sqrt{d} \cdot d^{3/2} \ll c\sqrt{d}n$   
 as  $d \leq n^{2/3}$  for  $c$  sufficiently small.

(ii) If  $n^{2/3} \leq d \leq 1/2 n$ , set  $\alpha = c/n/d^2$ .

Then  $\alpha \cdot \sqrt{d} = c/n/d^{3/2} \leq c \leq 1$  for  $c$  sufficiently small.

Also,  $\alpha nd = c \cdot n^2/d$  and  $d^3 \alpha^2 = c^2 \cdot n^2/d$ , so  
 choosing  $c$  small enough we get  $\text{disc}^1(G) = \Omega(n^2/d)$   $\square$ .