

Extremal Combinatorics: If  $G$  has minimum degree  $\delta \geq f(n)$ , then  $G$  contains  $H$  as a subgraph.

Robust Threshold: For which  $p$  can we still guarantee that  $G_p$  contains  $H$ ?

$G_p$  = random subgraph of  $G$  where each edge is retained with probability  $p$ .

Perturbed Random Graph: Let  $G$  be a graph with minimum degree  $\delta \geq \alpha n$ .

Let  $G' = G \cup G(n, p)$ . For which  $p$  can we guarantee that  $G'$  contains  $H$ ?

Eg.  $H$  = Hamiltonian cycle.

•  $p_c(H) = \Theta\left(\frac{\log n}{n}\right)$ . (Pósa)

• (Dirac) If  $G$  has minimum degree  $\geq \frac{n}{2}$  then  $G$  has a Hamiltonian cycle (tight).

(Krivelevich - Lee - Sudakov)

• Given  $G$  with minimum degree  $\geq \frac{n}{2}$ ,  $G_{c \log n / n}$  has a Ham. cycle with high probability.

• For  $\alpha \in (0, \frac{1}{2})$ : If  $G$  has minimum degree  $\delta \geq \alpha n$ , then for  $p = \frac{1}{n}$ ,  $G' = G \cup G(n, p)$  contains a Ham. cycle. (Bollobás - Frieze - Martin 03)

• If  $\delta(G) \geq \frac{n}{2}$ , then  $G$  contains at least  $(cn)^n$  Hamiltonian cycles

(Sárközy - Selkowitz - Szemerédi, constant by Cuckler and Kahn)

## Extremal Combinatorics

Thm (Corrádi - Hajnal): If  $\delta(G) \geq \frac{2}{3}n$ , then  $G$  contains a triangle factor.

Thm (Hajnal - Szemerédi): If  $\delta(G) \geq \frac{k}{k+1}n$ , then  $G$  contains a  $k_{k+1}$ -factor.

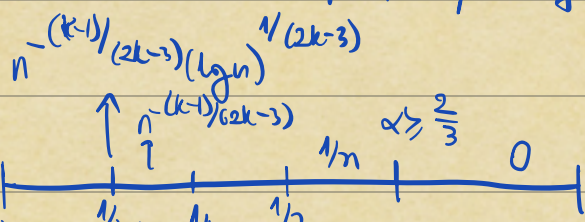
Dirac threshold  $\delta_{l,k} = \lim_{n \rightarrow \infty} \frac{t(n,k,l)}{\binom{n}{k-l}}$ ,  $t(n,k,l) = \min d$  s.t.  $\delta_2(H) \geq d \Rightarrow$  perfect matching.

Thm (Kondó - Sárközy - Szemerédi):

If  $\delta(G) \geq (\frac{1}{2} + \epsilon)n$ , then  $G$  contains a bounded degree spanning tree  $T$ .

Thm (KSS): If  $\delta(G) \geq (1 - \frac{1}{k+1})n$ , then  $G$  contains a  $k^{\text{th}}$  power of a Hamiltonian cycle.

Perturbed Random Graphs, Square of Ham. cycle



## Probabilistic Combinatorics

Robust version (ABCDJMS)

If  $\delta(G) \geq \frac{2}{3}n$ , then  $G_{C(\lg n)^{1/2} - \epsilon}$  contains a triangle factor.

(P.-Sah-Sawhney-Simkin):  $G_{C(\lg n)^{\frac{1}{k+1}} - \epsilon}$  contains a  $k_{k+1}$ -factor.

(PSS): If  $\delta(G) \geq \frac{k}{k+1}n$ , then  $G$  contains  $\binom{cn}{k+1}^{kn/(k+1)}$   $k_{k+1}$ -factors.

(PSSS, KKKOP): If  $\delta_2(H) \geq (\delta_{2,k} + \epsilon) \binom{n}{k-l}$  then  $H_{C(\lg n)^{1/k}}$  contains a perfect matching.

(PSS):

If  $\delta(G) \geq (\frac{1}{2} + \epsilon)n$ , then  $G_{C(\lg n)^{1/k}}$  contains a copy of  $T$ .

If  $\delta(G) \geq (1 - \frac{1}{k+1})n$ , then  $G_{Cn^{-1/k}}$  contains a  $k^{\text{th}}$ -power of Ham cycle ( $k \geq 2$ ). (KNP for  $k=2$ )

Thm (Nešetřil-P. 2c): For  $k \geq 3$ ,

$\delta(G) \geq (\frac{1}{k+1} + \epsilon)n$ , then  $GUG(n,p)$  contains a  $k^{\text{th}}$  power of the cycle for

$\frac{1}{\sqrt{n}}$   $\frac{1}{k}$   $\frac{1}{2}$   
 (Antonik, Dudek, Reiher, Ruciński, Schacht, Böttcher, Łuczak, Squeglia, Skokan)

$$p \geq Cn^{-1/(k-1)}$$

## Regularity Method:

Def:  $G = (A, B, E)$  is  $\varepsilon$ -regular if  $\forall A' \subseteq A, B' \subseteq B, |A'| \geq \varepsilon|A|, |B'| \geq \varepsilon|B|$ :

$$|d(A', B') - d(A, B)| \leq \varepsilon.$$

Here  $d(A', B') = \frac{|E(A', B')|}{|A'| |B'|}$ .

Thm (Szemerédi's Regularity Lemma):  $\forall \varepsilon > 0, \exists M(\varepsilon)$  st. the following holds. For any graph  $G$ , there is an <sup>equitable</sup> partition of  $V(G)$  into at most  $M(\varepsilon)$  parts  $V_1, \dots, V_m$  such that for all but at most  $\varepsilon m^2$  pairs of parts  $V_i, V_j$ ,  $G[V_i, V_j]$  is  $\varepsilon$ -regular.

„ $(A, B, E), |A| = |B| = n$ “

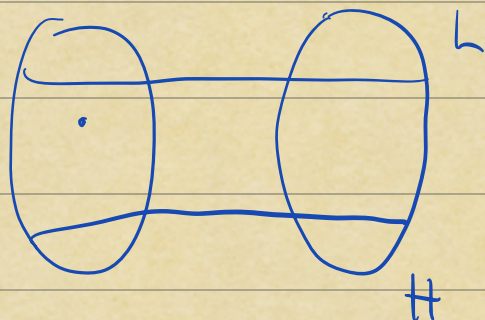
Def:  $G$  is  $\varepsilon$ -super-regular if  $G$  is  $\varepsilon$ -regular with density  $d$ , and every vertex has degree  $(d \pm \varepsilon)n$ .

We say  $G$  is  $(\varepsilon, \delta)$ -super-regular if  $G$  is  $\varepsilon$ -regular, and every vertex has degree at least  $\delta n$ .

Thm: For a bipartite graph  $G = (A, B, E)$  which is  $(\varepsilon, \delta)$ -super-regular, for any  $\bar{d} \leq \delta - C\varepsilon$ ,  $G$  contains a spanning subgraph which is  $\varepsilon'$ -super-regular with density  $\bar{d}$  where  $\varepsilon' = O(\varepsilon^{1/2})$ .

Pf: Let  $d_0$  be the density of  $G$ . By  $\varepsilon$ -regularity, all but  $O(\varepsilon n)$  vtx have degree  $(d_0 \pm \varepsilon)n$ .

Let  $\hat{G}$  be obtained by keeping each edge with prob  $\frac{d}{d_0}$ ,  $d = \bar{d} + C\varepsilon$ ,  
 let  $L, H$  be the vtx with degree less than  $(d - 2\varepsilon)n$  or higher than  $(d + 2\varepsilon)n$ .



For every  $v \in L$ , we add  $dn - \deg_{\hat{G}}(v)$  edges between  $v$  and  $(L \cup H)^c$ .

For  $v \in H$ , remove  $\deg_{\hat{G}}(v) - dn$  edges between  $v$  and  $(L \cup H)^c$ .

Now every vtx has degree  $(d \pm O(\varepsilon))n$ ,

and the density  $\geq \bar{d}$ .

Finally, obtain  $\bar{G}$  by removing  $\leq O(\varepsilon n)$  edges around each vtx s.t. edge density is exactly  $\bar{d}$ .

## The Blow-up Lemma

Then (Könlős - Sárközy - Szemerédi 97):  $\forall r, \Delta, \delta > 0, \exists \varepsilon > 0$  s.t.

Let  $R$  be a graph on  $r$  vtx. Consider a graph  $G$  on  $V_1 \cup \dots \cup V_r, |V_i| = n$ ,

s.t.  $(V_i, V_j)$  is  $(\varepsilon, \delta)$ -super-regular  $\forall \{i, j\} \in E(R)$ .

Let  $H$  be a graph with a homomorphism  $h: H \rightarrow R$  and  $|h^{-1}(i)| \leq n \forall i$ .

Assume  $\Delta(H) \leq \Delta$ . Then  $G$  contains a copy of  $H$ .

## Random Blow-up Lemma:

Def: Let  $\lambda$  be a distribution on  $\phi: X \rightarrow Y$ . We say  $\lambda$  is  $p$ -vertex-spread if  $\forall x_1, \dots, x_k, y_1, \dots, y_k$ :

$$\mathbb{P}_{\phi \sim \lambda}(\phi(x_i) = y_i \forall i) \leq p^k.$$

Thm (Neradoš - P. 24):  $\forall r, \Delta, \delta, \alpha > 0, \exists \epsilon, \beta > 0$  s.t.

Let  $R$  be a graph on  $r$  vtx. Consider a graph  $G$  on  $V_1 \cup \dots \cup V_r, |V_i| = n$ , s.t.  $(V_i, V_j)$  is  $(\epsilon, \delta)$ -super-regular  $\forall \{i, j\} \in E(R)$ .

Let  $H$  be a graph with a homomorphism  $h: H \rightarrow R$  and  $|h^{-1}(i)| \leq n \forall i$ .

Assume  $\Delta(H) \leq \Delta$ .

Let  $W \subseteq V(H), |W| \leq \beta n$  be given. For each  $x \in W$ , given  $W_x \subseteq V_h(G)$  of size  $\geq \alpha n$ .

Then there exists an  $O(1/n)$ -vertex-spread distribution  $\lambda$  on embeddings

$$\phi: H \rightarrow G \text{ s.t. } \phi(x) \in W_x \forall x \in W.$$

# Sketch of the Random Blow-up Lemma:

Key tool:

Lem (PSSS):  $\forall d, \delta > 0, \exists \epsilon > 0$ :  $G$   $(\epsilon, d)$ -super-regular then  
 $\exists \mu$  on perfect matchings in  $G$  which is  $O_\delta(V_n)$ -spread.

Let  $H_i = h_i^{-1}(i)$ . Choose  $B_i, D_i \subseteq H_i \setminus W$ :  $|B_i| = \delta_i n$ ,  $|D_i| = \beta_i n$ ,  
 BUD ind set  $(B = \cup B_i)$ , every 2 vtx of  $B \cup D$  have distance  $\geq 4$ .

Quasirandom embedding:

Def: For  $S \subseteq V(H)$ ,  $\phi: H[S] \rightarrow G$  quasirandom if

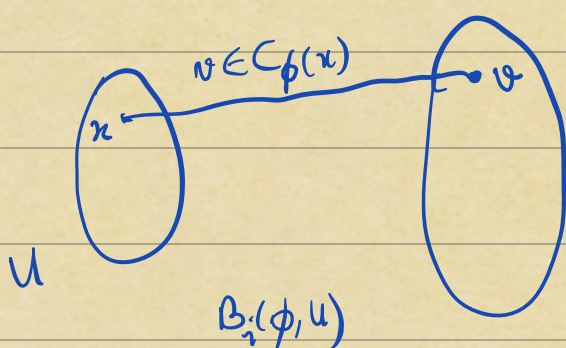
•  $\forall x \in V(H) \setminus S$ ,  $C_\phi(x) = W_x \cap \left( \bigcap_{y \in N_H(x) \cap S} N_G(\phi(y)) \right)$

has  $|C_\phi(x)| \geq (d - \epsilon)^{|N_H(x) \cap S|} |W_x|$ .

• for each  $i$ , at most  $\epsilon' |S| n$  pairs  $(x, y) \in (H_i \setminus (N_H(D) \cup S))^2$

has  $|C_\phi(x) \cap C_\phi(y)| \leq (d + \epsilon)^{|N_H(x) \cap S| + |N_H(y) \cap S|} n$ .

Lem:  $\phi$   $S$ -quasirandom. For  $i \in R$ ,  $U \subseteq H_i \setminus S$ :



If  $U \cap (W \cup N_H(D)) = \emptyset$ ,

$|U| \geq \delta_3 n$ ,

$|N_H(x) \cap S| = \ell \quad \forall x \in U$ , then

$B_i(\phi, U)$  is  $\epsilon''$ -regular, dens  $\geq (d - \epsilon)^\ell$ .

[Quasirandomness  $\Leftrightarrow \sum_{x,y} |C_\phi(x) \cap C_\phi(y)|^2$  small].

Embedding: Order  $V(H)$  s.t.  $N_H(b)$  comes first,  $b$  comes last.

Embed  $\{x_1, x_2, \dots, x_j\} = X_j$ .

Phase 1: For  $j \geq 0$ ,

• def  $A_{j+1} \subseteq C_\phi(x_{j+1}) \setminus \phi(X_j)$  the set of  $v$  s.t.

$$|N_G(v) \cap (C_\phi(y) \setminus \phi(X_j))| \geq (d-\epsilon) |C_\phi(y) \setminus \phi(X_j)|$$

$$\forall y \in N_H(x_{j+1}) \setminus X_j.$$

• Extending  $\phi$  to  $\phi(x_{j+1}) = v$  gives a quasirandom embedding.

Pick  $\phi(x_{j+1})$  from  $A_{j+1}$  unif at random.

• If  $j: s$ ,  $s = \delta_2 n$ , def  $L_j = \{v \notin X_j : |C_\phi(v) \setminus \phi(X_j)| < \delta_1 n\}$ ,  
move  $L_j$  forward (may include vtx of  $B$ ).

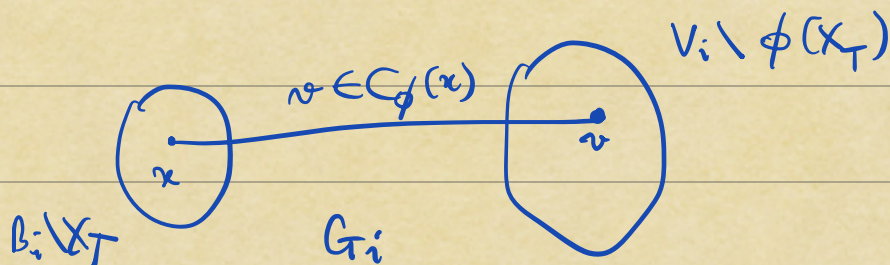
↳ All of  $N_H(b)$  already embedded

(Exceptional step) For  $j = |N_H(b)|$ ,  $E_i = \{v : |\{b \in B_i, v \in C_\phi(b)\}| < \delta_n |B_i|\}$

Take random injection  $\rho_i: E_i \rightarrow D_i$ , move  $\cup \rho_i(E_i)$  to beginning of ordering. For  $x \in \rho_i(E_i)$ , set  $\phi(x) = \rho_i^{-1}(x)$ .

Let  $T$  be the # vtx embedded.

Phase 2: For each  $i$ , consider



Sample  $O(1/n)$ -spread perfect matchings in  $G_i \rightarrow$  Ford embedding  $\phi$ .

Obs: From quasirandomness of embedding,  $L_j = \phi$  for  $j < |N_{\#}(B)|$ , so first embed all of  $N_{\#}(B)$ .

By lem,  $B_i(\phi, B_i)$   $\varepsilon''$ -regular, density  $\geq (d-\varepsilon)^\Delta \Rightarrow |E_i| \leq \varepsilon'' n < |D_i|$

Claim: For  $j$ 's,  $|L_j \setminus (W \cup N_{\#}(D))| < r(\Delta+1) \delta_3 n$ :

Choose  $l$  st.  $U_i^l \subseteq (\{i\} \cup L_j) \cup \{x \text{ s.t. } |N_{\#}(x) \cap X_j| = l \text{ has size } \geq \delta_3 n\}$   
 $\setminus (W \cup N_{\#}(D))$

lem  $\Rightarrow B_i(\phi, U_i^l)$   $\varepsilon''$ -regular.

#  $v_x$  moved up  $\leq O(\delta_3/\delta_2)n \ll \delta_0 n$  so  $\geq \frac{\delta_0 n}{2}$   $v_x$  of  $B$

unembedded.  $\Rightarrow \exists x \in U_i^l$  with  $> (d-\varepsilon)^l - \varepsilon'' n$  neighbors in unembedded

$v_x$ .  $\square$

$\Rightarrow$   $v_x$  moved up are always embedded before the next multiple of  $s$  phase.

Claim:  $\forall j, x \in V(H) \setminus X_j, |C_\phi(x) \setminus \phi(X_j)| \geq (d-\varepsilon)^{\Delta+1} \delta_1 n - 2s$ .

If  $x \notin L_j, |C_\phi(x) \setminus \phi(X_{j+s})| \geq (d-\varepsilon)^{d_1} \delta_1 n - s$ ,

$d_1 = |N_{\#}(x) \cap \{x_{j+s}, \dots, x_{j+2s}\}|$ .

If  $x$  embedded  $\rightarrow \checkmark$ , else consider  $x \in L_{j+s}$ .

$|C_\phi(x) \setminus \phi(X_{j+2s})| \geq (d-\varepsilon)^{d_2} ((d-\varepsilon)^{d_1} \delta_1 n - s) - s$ ,

$d_2 = |N_{\#}(x) \cap \{x_{j+s+1}, \dots, x_{j+2s}\}|$ .

$x$  will be embedded before step  $j+2s$ .  $\square$



Claim:  $|A_{j+1}| \gg \delta_2 n$ .

By regularity, all but  $O(\epsilon)n$   $v \in W_{x_{j+1}}$  has

$$|N_G(v) \cap C_\phi(y)| \geq (d-\epsilon) |C_\phi(y)|,$$

$$|N_G(v) \cap (C_\phi(y) \setminus \phi(x_j))| \geq (d-\epsilon) |C_\phi(y) \setminus \phi(x_j)| \quad \forall y \in N_+(x_{j+1}) \setminus X_j.$$

Let  $P = \{ (x,y) \in (H_i \setminus (N_+(x) \cup S))^2 : |C_\phi(x) \cap C_\phi(y)| \leq (d+\epsilon) \frac{|N_+(x) \cap S| + |N_+(y) \cap S|}{n}, x \text{ or } y \in N_+(x_{j+1}) \}$ .

For random  $v \in V_{H(x_{j+1})}$ ,  $\mathbb{E}[\# \text{pairs failing (2) after } \phi(x_{j+1})=v] \leq \epsilon P$ .

W.p.  $1-\epsilon'$ , at most  $\epsilon' n$  new pairs fail (\*).

$\Rightarrow$  All but  $O(\epsilon' n)$  vtx of  $C_\phi(x_{j+1}) \setminus X_j$  belongs to  $A_{j+1}$ .  $\square$

Phase 2:  $|b_i \setminus X_T| \geq |b_i| - O(\delta_3/\delta_2)n$ ,

By Lem,  $B_i(\phi, b_i \setminus \phi(X_T))$   $\epsilon''$ -regular.

$F_i = V_i \setminus \phi(X_T)$ : From exceptional step, every  $v \in F_i$  has deg

$$\geq \delta_1 |b_i| - O(\delta_3 n / \delta_2),$$

while Claim  $\Rightarrow |C_\phi(b) \setminus \phi(X_T)| \geq \delta_2 n$ .

$\} \Rightarrow \exists$  spread perfect matching.