



# Lecture 7

## § Stability method

Tackle an extremal problem

- Step 1 : Obtain asymp. result ;
  - Step 2 : Obtain Stability result ;
  - Step 3 : Use the structural info. from the Stability

Idea → to slowly remove imperfection to get an exact result.

Illustration: Thm  $\text{ex}(n, C_{2k+1}) = \lfloor \frac{n^2}{4} \rfloor$

( Recall : Special case of color-critical gr }  $\stackrel{=}{\text{ex}}(n; k_3)$

$$\text{Step 1 E-S-S : } \text{ex}(n, C_{k+1}) = \frac{n^2}{4} + o(n^2)$$

Step 2: E-Sim:

Lem  $\forall \alpha < \omega_1 \exists n \in \omega$  s.t.  $\text{TFH}(\Lambda) \cap \omega = n$

$$\forall G \quad n - vx \left\{ \begin{array}{l} C_{2k+1} \text{-free} \\ e(G) \geq \frac{n^2}{4} - \alpha n^2 \end{array} \right. \Rightarrow \begin{array}{l} G \text{ can be made} \\ \text{bipartite by removing} \\ \leq 3\alpha n^2 \text{ edges} \end{array}$$

Idea : • Take max-cut  $X \cup Y$

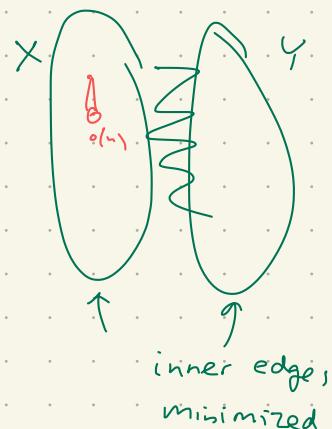
$$\underline{\text{STS}} \quad e_G(x) + e_G(y) = 0$$

$$\therefore |X| = |Y| \approx n_2$$

- (Weaker)  $\max \text{inner deg} = o(n)$

a single

• Not inner edge allowed



Pf • Let  $G_{\text{ex}}$  be an  $N$ -vx extremal  $C_{2k+1}$ -free gr

$$\text{Lower bd: } e(G_{\text{ex}}) \geq \lfloor \frac{N^2}{4} \rfloor = \binom{\lfloor \frac{N}{2} \rfloor}{2} = O(2^N)$$

- boost min deg: find  $G \subset G_{\text{ex}}$  s.t.  $\delta(G) \geq (\frac{1}{2} - \varepsilon)|G|$

Start w/  $G_0 = G_{\text{ex}}$

Do: • if  $\exists v_0 \in V(G_0)$  w.  $d_{G_0}(v_0) < (\frac{1}{2} - \varepsilon)N$

let  $G_1 = G_0 - v_0$

At  $i$ th step

$\exists v_{i-1} \in V(G_{i-1})$  w.  $d_{G_{i-1}}(v_{i-1}) < (\frac{1}{2} - \varepsilon)(N-(i-1))$

let  $G_i = G_{i-1} - v_{i-1}$

Claim this process must terminate w/ a desired

gr  $G_k = G$

$$\underline{\text{Pf}} \quad e(G_k) \geq \frac{N^2}{4} - \sum_{i=1}^k (\frac{1}{2} - \varepsilon)(N+i-i)$$

$$\geq \frac{N^2}{4} - (\frac{1}{2} - \varepsilon)Nk + (\frac{1}{2} - \varepsilon) \sum_{i=1}^k (i-1)$$

Write  $n = |G_k| = N - k$

$$= \frac{(n+k)^2}{4} - (\frac{1}{2} - \varepsilon)(n+k)k + \dots$$

$$= \frac{n^2}{4} + \varepsilon nk + \frac{\varepsilon}{2}k^2 + O(k)$$

When  $k = \mathcal{O}(N) = \mathcal{O}(n) \Rightarrow \frac{n^2}{4} + \mathcal{O}(n^2)$

↳ E-S-S.

- We have  $n - ux \leq G \subseteq G_{ex}$   $\left\{ \begin{array}{l} C_{2k+1}-\text{free} \\ \delta(G) \geq (\frac{1}{2} - \varepsilon)n \end{array} \right.$
- Consider Max-cut  $X \cup Y$  for  $G$   $\Rightarrow e(G) \geq \frac{n^2}{4} - \frac{\varepsilon}{2}n^2$

Apply Stability on  $G$

$$\Rightarrow e_G(X) + e_G(Y) \leq \frac{3\varepsilon}{2}n^2$$

Claim  $|X| = |Y| = \frac{n}{2} \pm 2\sqrt{\varepsilon}n$

Pf O.w.

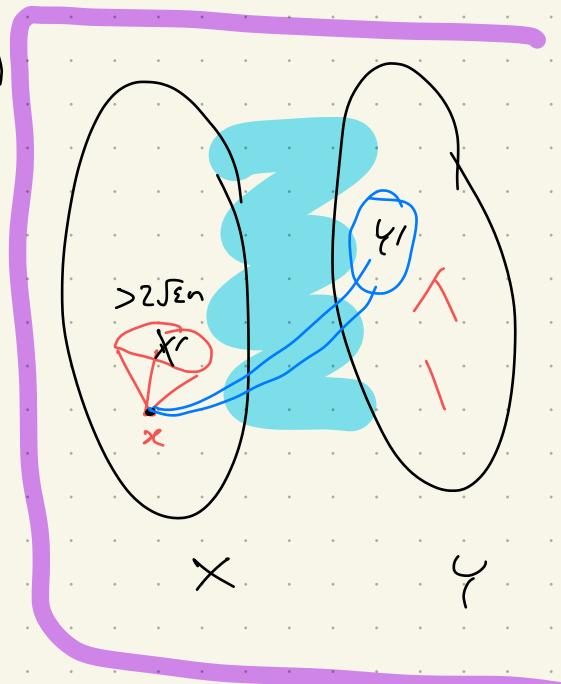
$$e(G) \leq |X||Y| + e_G(X) + e_G(Y)$$

$$\leq \frac{n^2}{4} - 4\varepsilon n^2 + \frac{3\varepsilon}{2}n^2$$

$$< \frac{n^2}{4} - \frac{\varepsilon}{2}n^2$$

Claim  $\Delta(G[X]), \Delta(G[Y])$

$$\leq 2\sqrt{\varepsilon}n.$$



Pf Suppose not,  $\exists x \in X$

$$d(x, X) > 2\sqrt{\varepsilon}n$$

$$\text{Max-cut} \Rightarrow d(x, Y) \geq d(x, X) > 2\sqrt{\varepsilon}n$$

$$\text{Let } N(x, X) = X', \quad N(x, Y) = Y'$$

As  $G$  is  $C_{2k+1}$ -free, the bip. gr

$$\Rightarrow G[X', Y'] \text{ is } P_{2k}\text{-free}$$

Thm (Erdős-Gallai)  $\text{ex}(n, P_t) \leq \frac{t-2}{2} \cdot n$

Rank optimal when  $t-1/n$ :

$$e_G(X) + e_G(Y) \quad \begin{matrix} \downarrow \\ \text{cross edges} \end{matrix}$$

$$e(G) \leq \frac{3\epsilon n^2}{2} + (|X||Y| - (|X'||Y'| - kn))$$

$\leftarrow \frac{n}{t-1}$  many

$$\leq \frac{3\epsilon n^2}{2} + \frac{n^2}{4} + kn - 4\epsilon n^2$$

$$< \frac{n^2}{4} - \frac{\epsilon}{2} n^2$$

$\nwarrow$  # missing edges  
between  $X'$  &  $Y'$

$$G \left\{ \begin{array}{l} \cdot |X| = |Y| = \frac{n}{2} \pm 2\sqrt{\epsilon} n \\ \cdot \delta(G) \geq (\frac{1}{2} - \epsilon) n \\ \cdot \Delta(G[X]), \Delta(G[Y]) \leq 2\sqrt{\epsilon} n \end{array} \right.$$

We now show no edge is  $X$  or  $Y$ .

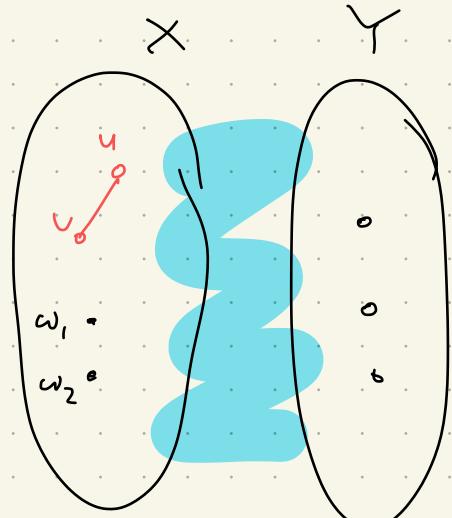
Suppose  $uv$  in  $X$  form an edge

$\forall$  vertex  $x \in X$ , almost full deg to  $Y$ :

$$d(x, Y) \geq \delta(G) - \Delta(G[X])$$

$$\geq (\frac{1}{2} - \epsilon) n - 2\sqrt{\epsilon} n$$

$$\geq (1 - 10\sqrt{\epsilon}) |Y|$$



$k=3$   
 $C_7$ -free

Pick arbitrary  $w_1, \dots, w_{k-1} \in X - u - v$

$\Rightarrow u, v, w, \dots, \omega_{k-1}$  have  $\geq (1 - (k+1)\sqrt{\varepsilon})|Y| \geq 3$

$\Rightarrow C_{k+1} \subseteq G \quad \square$

So  $G$  is bipartite w/ partition  $X \cup Y$

$$e(G_{\text{ex}}) \leq e(G) + \sum_{i=n+1}^N \left(\frac{1}{2} - \varepsilon\right) i \\ < \frac{N^2}{4} \quad \text{if } n < N$$

$\Rightarrow G_{\text{ex}}$  must be  $G \Rightarrow G_{\text{ex}}$  bip

$\Rightarrow G_{\text{ex}} \cong T_{N,2} \quad \square$

HW

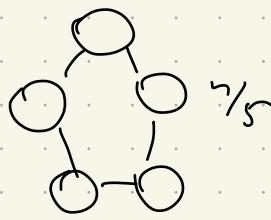
sent to

[zhuo.wu@warwick.ac.uk](mailto:zhuo.wu@warwick.ac.uk)

• If  $n = ux$   $\Delta$ -free  $G$

if  $\delta(G) > \frac{2n}{5} \Rightarrow G$  bipartite

Optimal



$$\delta(G) = \frac{2n}{5}$$

non-bip  
 $\Delta$ -free

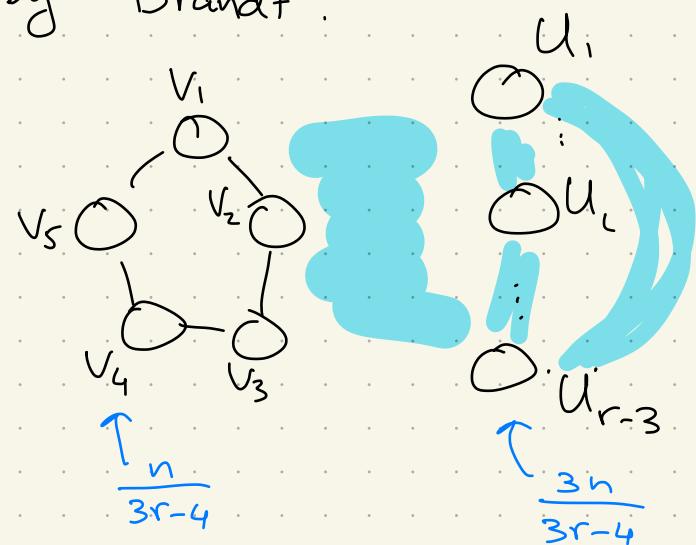
# Thm Andrasfai - Erdős - Sós

$\forall n \text{-ux } K_r\text{-free gr } G$

$$\delta(G) > \frac{3r-7}{3r-4} \cdot n \Rightarrow G \text{ is } (r-1)\text{-partite.}$$

We give a short pf by Brandt.

Rmk Bound is tight:

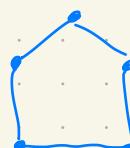


Def (5-wheel-like graph  $W_{r,k}$ )

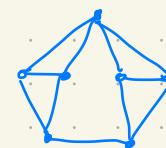
$$0 \leq k \leq r-3$$

- top ux  $v$
- bottom edge  $w_1, w_2$
- two  $(r-2)$ -cliques  $Q_1, Q_2$   
with  $|Q_1 \cap Q_2| = k$
- $v \sim Q_1 \& Q_2$
- $w_1 \sim Q_1$
- $w_2 \sim Q_2$

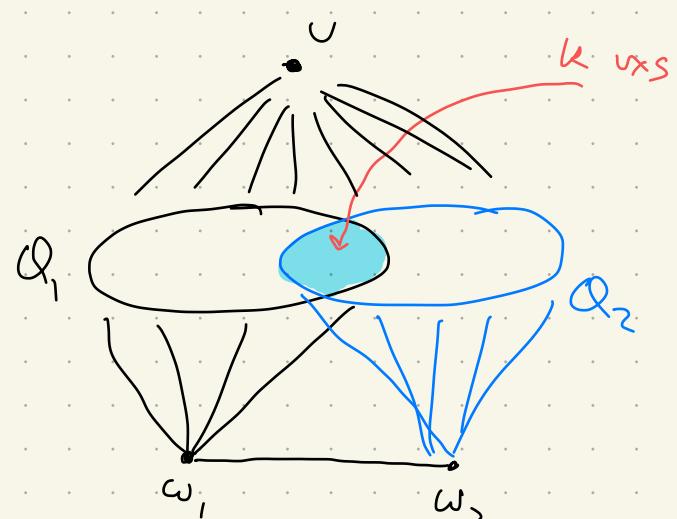
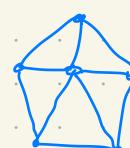
Ex  $W_{3,0}$



$W_{4,0}$



$W_{4,1}$



The nice idea of Brandt is the follow fact generalizing the fact that

"every nonbipartite maximal  $\Delta$ -free gr contains a 5-cycle"

Lem  $\forall G$  • maximal  $K_r$ -free  $\Rightarrow G$  contains a  
• non- $(r-1)$ -partite 5-wheel-like gr

Pf: • If  $G$  is complete multipartite  
 $\Rightarrow$  # parts  $\leq r-1$  as  $G$  is  $K_r$ -free  $\hookrightarrow$

Thus  $G$  is not complete multipartite

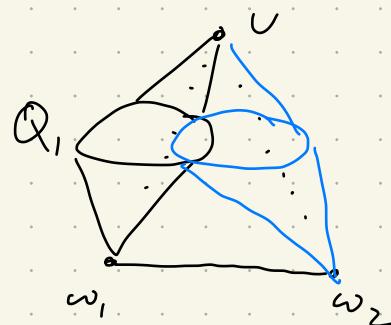
$\Rightarrow$  non-adjacency is NOT an equivalence relation

$\Rightarrow G$  contains

$$vw_1 \notin E$$

$$vw_2 \notin E$$

$$w_1 w_2 \in E$$



•  $G$  is maximal  $K_r$ -free  $\forall i \in [2]$

$\Rightarrow N(v) \cap N(w_i)$  contains a copy  $Q_i$   
of  $K_{r-2}$

for otherwise  $G + vw_i$  is still  $K_r$ -free  $\hookrightarrow$  maximality of  $G$ .

PF (AES:  $\delta(G) > \frac{3r-7}{3r-4}n$ ,  $\wedge$  Kr-free  $\Rightarrow$  (r-1)-partite)

- By adding edges if necessary, we may assume  $G$  is maximal Kr-free.

Lem  $\Rightarrow G$  contains  $W_{r,k}$

- Pick a copy  $W$  of  $W_{r,k}$  with maximum value  $k$ .

Define  $X = \text{common neighbor of } Q_1 \cap Q_2$

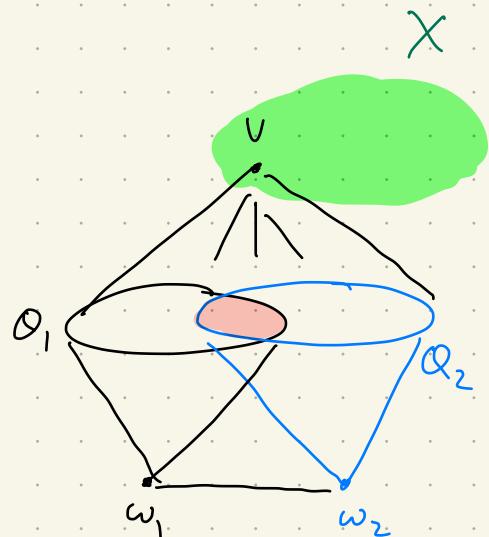
We will bound  $|X|$  in two ways.

- i) By double counting deg sum in  $W$ .

$\begin{cases} G \text{ is Kr-free} \\ W \text{ contains } \cong K_{r-1} \end{cases} \Rightarrow \forall u \in V \setminus W, u \text{ is not completely joined to } W$

Write  $p = |W|$

$$\Rightarrow d(v, W) \leq p-1$$



Claim  $\forall x \in X, d(x, W) \leq p-3$

PF:  $G$  is Kr-free

$\Rightarrow x$  has a non-neighbor in  $Q_i \cup v$  and  $Q_i \cup \{w_i\} \quad \forall i \in [2]$

Case 1  $x \sim v$

By the above obs.

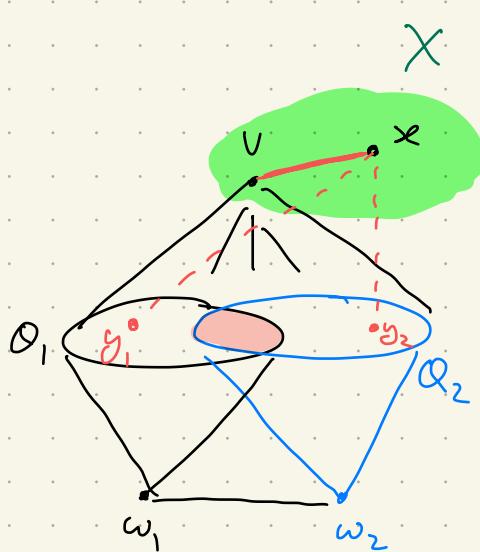
$\Rightarrow x$  has a non-neighbr

$$y_1 \in Q_1 \setminus Q_2, y_2 \in Q_2 \setminus Q_1$$

Consider  $v, w_1, w_2, Q'_1, Q'_2$

$$\text{where } Q'_i = Q_i - y_i + x$$

$\Rightarrow$  we get  $W_{r, k+1}$



$\hookrightarrow$  maximality of  $W$ .

$\Rightarrow x \not\sim v \Rightarrow d(x, W) \leq p-3$  😊

&  $x$  has non-neighbr  
in (disjoint)  $Q'_i \cup w_i \setminus (Q_1 \cap Q_2)$

• deg sum in  $W$  deg sum of vs to  $W$

$$\textcircled{1} \dots p \cdot \delta(G) \leq \sum_{w \in W} d(w) \leq (p-3)|X| + (p-1)(n-|X|)$$

• On the other hand, by pigeonhole

as  $X = \text{common neighbor of } Q_1 \cap Q_2$

$$\Rightarrow |X| \geq k \delta(G) - (k-1)n \dots \textcircled{2}$$

$$\textcircled{1} + \textcircled{2} \Rightarrow \delta(G) \leq \frac{2r+k-4}{2r+k-1} n \stackrel{k \leq r-3}{\leq} \frac{3r-7}{3r-4} n. \hookrightarrow \textcircled{3}$$



AES:  $n$ -ux  $\Delta$ -free,  $\delta(G) > \frac{2n}{5} \Rightarrow \underline{\text{bip}}$

$$\chi(G) \leq 2$$

Q: How much can we lower the min-deg  
if we only require bdd chr. #?

This <sup>the</sup> chromatic threshold problem.

### Extensions

Def  $\text{ex}(n, T, H) = \max$  # copies of  $T$  in an  
 $n$ -ux  $H$ -free gr

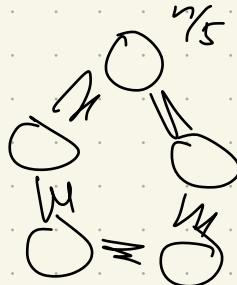
$$\text{ex}(n, H) = \text{ex}(n, K_2, H)$$

• Erdős - Zykov solve the clique case:

$H \subset t$   $\text{ex}(n, K_s, K_t)$  is maximized by  $T_{n, t-1}$

- Alon - Shikhelman '16
- Gishboliner - Shapira (cycles)
- Grzesik, Hatami - Hladky - Král' - Norin - Razborov

$$\text{ex}(n, S_5, K_3) \leq \left(\frac{n}{5}\right)^5$$



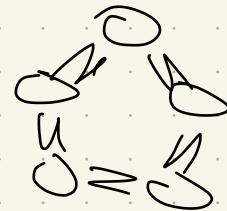
Two conj of Erdős (\$250)

Conj (Sparse halves)  $\forall n \in \mathbb{N}$   $\Delta$ -free gr  $G$

contains a set of  $\frac{n}{2}$  vxs inducing  $\leq \frac{n^2}{50}$  edges

Rmk

Best possible if true:



Conj  $\forall n \in \mathbb{N}$   $\Delta$ -free gr can be made

bipartite by removing  $\leq \frac{n^2}{25}$  edges