

The codegree Turán density

Def: Given a k -graph H and $S \subseteq V(H)$, the degree of S is $d(S) = |\{e \in E(H) : S \subseteq e\}|$. For $i \in [k-1]$, we set $\delta_i(H) = \min_{S \subseteq V(H)^{(i)}} d(S)$ and call it the minimum i -degree of H . For $i=k-1$, $\delta_{k-1}(H)$ is also called the minimum codegree of H .

Def: Given a k -graph F and $n \in \mathbb{N}$, we define $ex_{k-1}(n, F) = ex_{co}(n, F) = \max\{\delta_{k-1}(H) : H \text{ is an } F\text{-free } k\text{-graph on } n \text{ vertices}\}$

Def+Lemma (Mubayi, Zhao): The limit $\lim_{n \rightarrow \infty} \frac{ex_{k-1}(n, F)}{n}$ exists. It is called the codegree Turán density of F and we denote it by $\delta(F)$.

We omit the proof here.

Thm (Mubayi): $\delta(K_3) = \frac{1}{2}$

proof: For the lower bound, it is enough to note that the balanced "bipartite" 3-graph (two vertex classes with all those triples as edges which intersect both vertex classes) has minimum codegree $\sim \frac{n}{2}$ and does not contain K_3 .

For the upper bound, we use the following lemma.

Def: Let F_2 be the 3-graph obtained from $E_k K_4^{(3)}$ by duplicating the apex, i.e.

$$F_2 = (\{1, \dots, 5\}, \{134, 135, 145, 234, 235, 245\}).$$

Set $\mathcal{F}_2 = \{F_2\}$ and for $t \geq 2$, let \mathcal{F}_t be the family of 3-graphs obtained as follows. Take $F \in \mathcal{F}_{t-1}$, add two new vertices x and y and add t edges of the form axy , $a \in V(F)$.

lemma (Mubayi, Rödl): For $F \in \mathcal{F}_t$, we have

- 1) F has $2t+1$ vertices;
- 2) Every set of $t+2$ vertices in F spans at least one edge;
- 3) $\pi(\mathcal{F}_t) \leq \frac{1}{2}$ for $t \geq 2$.

Let us first use this lemma to prove the theorem.

proof of thm: Given $\varepsilon > 0$, choose $n_0, t \in \mathbb{N}$ with $n_0^{-1} \ll t^{-1} \ll \varepsilon$. Let H be a 3-graph on $n \geq n_0$ vertices with $\delta_{k-1}(H) \geq (\frac{1}{2} + \varepsilon)n$. Recall that by supersaturation, we have $\pi(\mathcal{F}_t) = \pi(\mathcal{F}'_t)$, where \mathcal{F}'_t is the family of 3-graphs obtained from \mathcal{F} by replacing each $F \in \mathcal{F}$ with $F(2)$. Further note that $e(H) \geq \frac{1}{3} \binom{n}{2} (\frac{1}{2} + \varepsilon)n > (\frac{1}{2} + \varepsilon) \binom{n}{3}$, whence we conclude that $F(2) \subseteq H$ for some $F \in \mathcal{F}_t$, using Part 3) of the lemma. Let $V(F(2)) = X \cup X'$ such that every vertex in F has one copy in X and one in X' . For every $x \in X$, denote its copy in X' by x' . Further, let $N(x, x') = \{z \in V(H) : xx'z \in E(H)\}$ be the coneighbourhood of x and x' . We then have

$$\sum_{x \in X} |N(x, x')| \geq (2t+1) (\frac{1}{2} + \varepsilon)n,$$

where we used Part 1) of the lemma and

the minimum codegree condition on H . Thus, there is some $y \in V(H)$ with $y \in N(x, x')$ for at least $(\frac{1}{2} + \epsilon)(2t+1) \geq t+2$ different choices of $x \in X$. By Part 2) of the lemma, there are $a, b, c \in X$ with $abc \in E(H)$ and $y \in N(x, x')$ for all $x \in \{a, b, c\}$. Now $H[\{a, b, c, a', b', c', y\}]$ contains a copy of \mathbb{F} .

let us now prove the lemma.

proof of lemma:

Part 1): This is an easy induction.

Part 2): We proceed by induction on t . For $t=2$, it is easy.

$t-1 \rightarrow t$: Let $F' \in \mathcal{F}_t$ be obtained by taking $F \in \mathcal{F}_{t-1}$, adding vertices x and y , and adding t edges axy with $a \in V(F)$. Now consider a set S of $t+2$ vertices. If $|S \cap V(F)| \geq t+1$, we get an edge in S by induction. Otherwise, $|S \cap V(F)| = t$ and $x, y \in S$, and since $|V(F)| = 2t-1$, there is some $a \in S$ with $axy \in E(F)$.

Part 3): We commence by recalling $\pi(K_4^{(3)}) \leq \frac{1}{3}$, which implies $\pi(\mathbb{F}_2) = \pi(\mathbb{F}_2) \leq \frac{1}{3}$ using supersaturation.

Hence, there is some constant $c > 2$ such that

$ex(n, \mathbb{F}_2) \leq \frac{1}{2} \binom{n}{3} + 2cn^2$ for all $n \in \mathbb{N}$. Now we

argue by induction on t that

$ex(n, \mathbb{F}_t) \leq \frac{1}{2} \binom{n}{3} + \epsilon cn^2$ for all $n \in \mathbb{N}$.

$t-1 \rightarrow t$: Use induction on n . For $n = 2t+1$, this is easy to see.

$n-1 \rightarrow n$:

Let H be a 3-graph on n vertices with $e(H) \geq \frac{1}{2} \binom{n}{3} + tcn^2$ edges. By the induction on t , we know that $F \subseteq H$ for some $F \in \mathcal{F}_t$. Let L be the multigraph on $V(H) \setminus V(F)$ whose edges are given by the edges in the links of the vertices in F . By definition of \mathcal{F}_t , if any edge in L has multiplicity $\geq t$, this pair together with F forms a copy of an element of \mathcal{F}_t in H . If this does not happen, there is a vertex $v \in V(L)$ with

$$d_H(v) \leq \underbrace{\frac{t-1}{2t-1} \binom{n-2t+1}{2}}_{\substack{1 \text{ vtx in } F, \text{ two} \\ \text{in } V(H) \setminus V(F)}} + \underbrace{(2t-2)(n-2t+1)}_{\substack{2 \text{ vtx in } F, \text{ one in} \\ V(H) \setminus V(F)}} + \underbrace{\binom{2t-1}{2}}_{\substack{3 \text{ vtx} \\ \text{in } F}}$$

$$< \frac{t-1}{2t-1} \binom{n-1}{2} + 2tn$$

Deleting v leaves us with a 3-graph on $n-1$ vertices with at least

$$e(H) - \frac{t-1}{2t-1} \binom{n-1}{2} - 2tn \geq \frac{1}{2} \binom{n}{3} + tcn^2 - \frac{1}{2} \binom{n-1}{2} - 2tn > \frac{1}{2} \binom{n-1}{3} + ct(n-1)^2$$

edges. By induction, this 3-graph contains an element of \mathcal{F}_t .