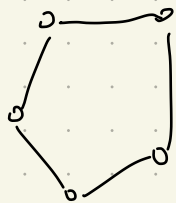




Lecture 1

• Graph (V, E)

vertices \uparrow edges (binary symm. on V)



5 vxs

5 edges

• Basic graph families

• paths :

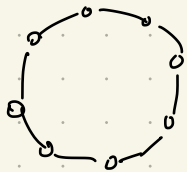


P_9

P_k

$k = \#$ vxs in path

• cycle :

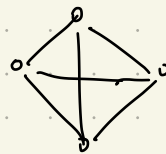


C_8

(order)

• complete graph

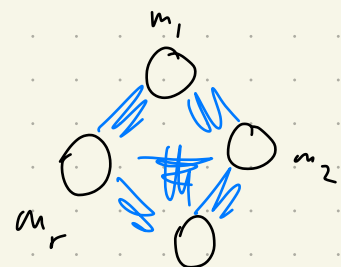
K_t



$t=4$

• complete multipartite graph

K_{m_1, m_2, \dots, m_r}

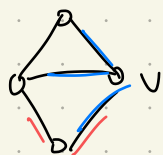


• r parts

• no edge in each part

• all edges between different parts

• degree of a vertex v : $d(v) = \#$ edges incident to v



$\deg = 3$

$u = \deg = 2$



群聊: Spring Course 2024



该二维码7天内(3月12日前)有效, 重新进入将更新

Prop $\forall G, \# \text{ edges in } G = \sum_{v \in V(G)} d(v)$

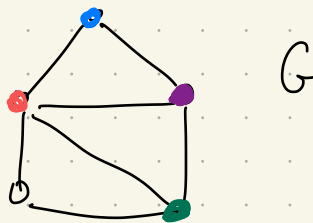
- $\delta(G)$ min deg of G
- $\Delta(G)$ max deg of G

Def Subgraph containment

Given two graphs H & G , we say H is a subgraph of G , denoted by $H \subseteq G$, if

\exists injective map $\varphi: V(H) \rightarrow V(G)$ preserving adjacencies
 i.e. $\forall uv \in E(H) \Rightarrow \varphi(u)\varphi(v) \in E(G)$

Ex



If no such map: G is H -free

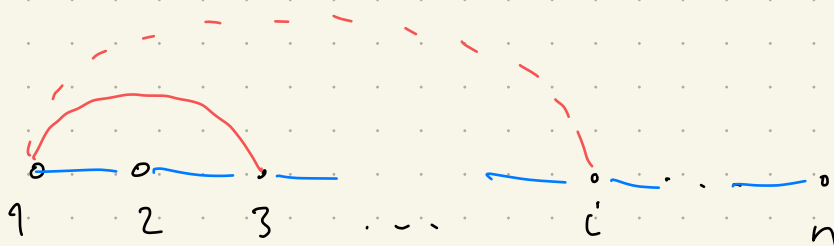
Prop $\forall G$ w/ average deg $d(G) = \frac{\sum_{v \in V(G)} d(v)}{|V(G)|}$

$d(G) \geq 2 \Rightarrow \exists \text{ cycle } \subseteq G$

Exer. Prove it \uparrow

Exer Is this bound tight / optimal? (YES)

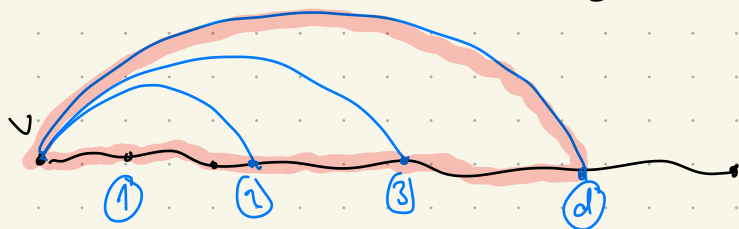
Rmk : We cannot say anything about this cycle
(no control)




$$d(G) = 2 \Leftrightarrow e(G) = n = |V(G)|$$

Prop $\forall G, \delta(G) \geq d \Rightarrow \exists$ cycle of length $\geq d+1$ in G

PF : • Take a longest path, say starting at some vertex v



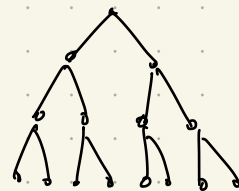
longest $\Rightarrow N(v)$

\leftarrow set of all neighbors of $v \Rightarrow$ desired cycle 

Rmk Bound optimal: Consider $G = K_{d+1}$

• (trees) A tree is a graph without any cycle (as a subgraph)
(acyclic)

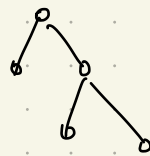
Ex



Ex \forall tree on t vertices T

$t=5$

$$\forall G \quad \delta(G) \geq t-1 \Rightarrow T \subseteq G$$



Rmk Optimal $G = K_t$

Mantel's thm

Thm (Mantel)

$\forall G$

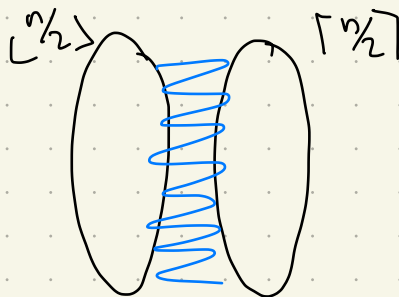
• Δ -free (K_3 -free)

• n -vertex

$$\Rightarrow e(G) \leq \lfloor \frac{n^2}{4} \rfloor$$

Rmk

Optimal!

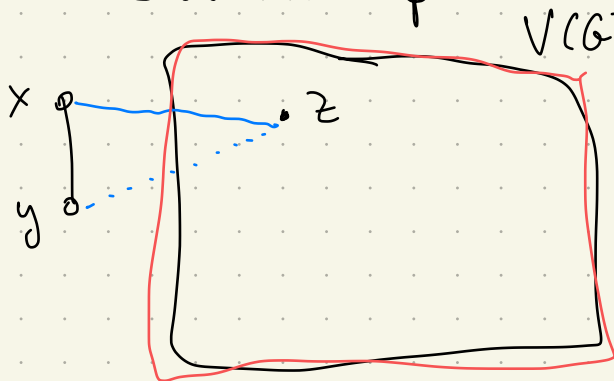


PF 1

Induction on the number of vertices n .

• Base case: easy.

• Inductive step:



$\forall u \in V(G) \setminus \{x, y\}$

is adjacent to ≤ 1

vertex in $\{x, y\}$

by Δ -freeness

$$e(G) \leq 1 + n - 2 + \left\lfloor \frac{(n-2)^2}{4} \right\rfloor \leq \left\lfloor \frac{n^2}{4} \right\rfloor.$$

↑
x, y edge

↑

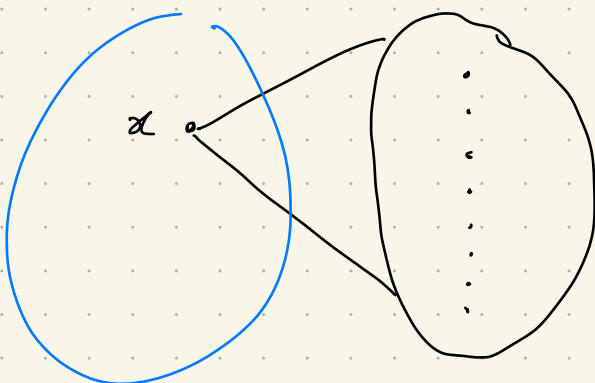
edges from $V \setminus \{x, y\}$ to $\{x, y\}$

↑ I.H.



Pf 2 (Max deg v_x)

Consider a vertex of maximum deg $d(v_x) = \Delta(G)$



$Y = V \setminus X$

$X = N(v)$

• Δ -free

\Rightarrow no edge in X

\Rightarrow every edge in G

is incident a vertex in Y

$$\Rightarrow e(G) \leq \sum_{y \in Y} d(y) \leq |X||Y| \leq \left(\frac{|X|+|Y|}{2} \right)^2 = \frac{n^2}{4}$$

↑
choice of x



$K_{\lfloor n/2 \rfloor, \lceil n/2 \rceil}$

Ex 1 Prove that $K_{\lfloor n/2 \rfloor, \lceil n/2 \rceil}$ is the unique maximizer (extremal graph).

— Sets —

- Odd town even town problem

Setup : • A town of n ppl

$$\{1, 2, 3, \dots, n\} =: [n]$$

- form into clubs C_1, C_2, \dots

$$C_i \subseteq [n]$$

Need \forall distinct clubs C_i, C_j

$$|C_i \cap C_j| \text{ even number}$$

Q1 If each club C_i admits even # ppl,
then max # clubs there can be?

A: We can have exponentially many.

Group ppl into pairs $\{1, 2\}, \{3, 4\}, \dots, \{n-1, n\}$ ^{even}

$$\# \text{ pairs} = n/2$$

Clubs = all possible combinations of pairs
 $2^{n/2}$

Q2 What if each club has odd size?

Try $C_i = \{i\}$ # clubs = n

$$|C_i \cap C_j| = 0 \text{ even}$$

Thm $\forall \mathcal{F} \subseteq 2^{[n]}$ ← power set of $[n]$: family of all subsets of $[n]$.

① $\forall F \in \mathcal{F}$, $|F|$ odd

② \forall distinct $F, F' \in \mathcal{F}$

$|F \cap F'|$ even

sets in \mathcal{F}



$$\Rightarrow |\mathcal{F}| \leq n$$

PF Consider $\mathbb{1}_F$ for each $F \in \mathcal{F}$

$n=5$

$$\{1, 2, 3, 4, 5\} = [n]$$

$$F = \{2, 4, 5\}$$

$$\mathbb{1}_F = (0, 1, 0, 1, 1)$$

① $\Rightarrow \langle \mathbb{1}_F, \mathbb{1}_F \rangle = 1$ (in \mathbb{F}_2)

② $\Rightarrow |F \cap F'| = \langle \mathbb{1}_F, \mathbb{1}_{F'} \rangle$

$$= 0 \text{ (in } \mathbb{F}_2)$$

$\Rightarrow \{ \mathbb{1}_F : F \in \mathcal{F} \}$ pairwise orthogonal

$$\subseteq \mathbb{F}_2^n$$

$$\Rightarrow |\mathcal{F}| = |\{ \mathbb{1}_F : F \in \mathcal{F} \}| \leq \dim \mathbb{F}_2^n = n$$



Ramsey

Philosophy: No complete disorder

\exists highly ordered substructure in any suff. large system.
↑
large

Consider two edge-coloring of complete graphs.
blue/red

Goal: find large monochromatic complete subgraphs

Def Ramsey number $R(s, t) = \min n \in \mathbb{N}$ s.t.
 $s, t \in \mathbb{N}$

\forall 2 edge-coloring of $K_n \implies \exists$ either a blue K_s
blue/red or a red K_t

Warm $R(2, 3) = 3$

Lower bound: by construction

$$R(2, 3) > 2$$

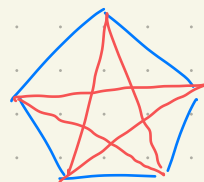


Upp bd

$$R(2, 3) \leq 3$$

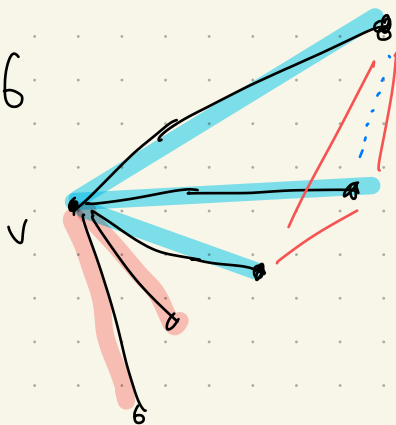


• $R(3,3) = 6$ • $R(3,3) > 5$



$R(3,3) \leq 6$

PF

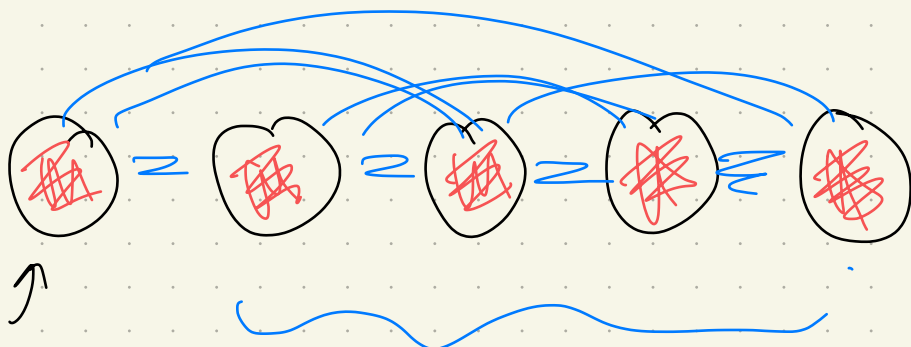


Pigeonhole $\Rightarrow \geq 3$ edges
Same color

Lower bound on

$R(k,k) > n$: Need to find a 2-coloring of K_n w./ no red/blue K_k .

Ex $R(k,k) > (k-1)^2$



↑
each gp size $k-1$

#gps = $k-1$

total # vcs
= $(k-1)^2$

• Use probabilistic method to prove lower bound on $R(k,k)$

(Implicit construction)

Thm (Erdős) $R(k, k) > n = (1+o(1)) \frac{k}{\sqrt{2e}} \cdot 2^{k/2}$

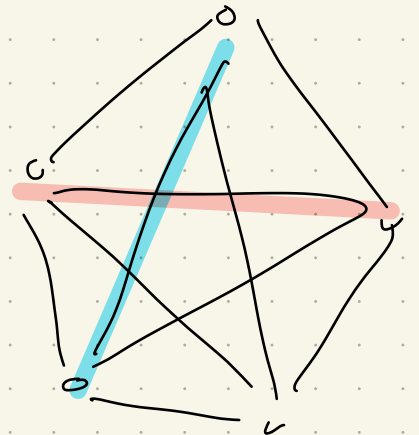
Pf: Consider a random edge-coloring of K_n

\forall edge $ij \in \binom{[n]}{2}$

\uparrow all pairs in $[n]$

Color ij $\begin{cases} \text{red} \\ \text{blue} \end{cases}$ w./ prob $\frac{1}{2}$

(independent of all other edges)



\mathbb{E} # monochromatic K_k

linearity of expectation

$$= \binom{n}{k} \cdot \Pr(\text{a fixed } k\text{-set is monochromatic})$$

$$= \binom{n}{k} \left(\frac{1}{2}\right)^{\binom{k}{2}} \cdot 2 < 1$$

$\Rightarrow \exists$ a 'good' coloring

