



Lecture 32

§ C_6 -free incidence graphs.

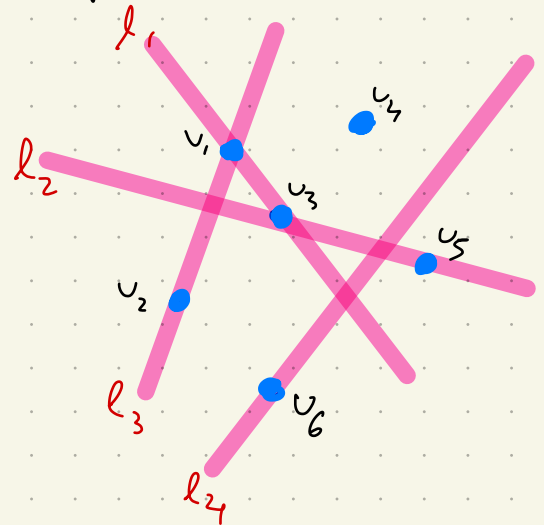
Given a P, L pts/lines arrangement on the plane $\mathbb{R}^2$,

an incidence is a pair $(p, l) \in P \times L$

s.t. $p \in l$.

Ex $|P| = 6, |L| = 4$

incidence = $I(P, L) = 7$

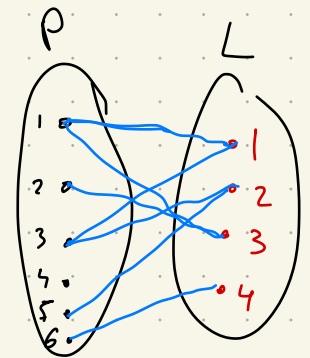


The pt/line incidence graph for an arrangement (P, L) is a bip. graph w/ parts P & L

and $pl \in G \iff p \in l$

• $|I(P, L)| = e(G)$

• $\deg(pt) = \# \text{ lines thru it.}$



Rmk pt/line incidence gr. is C_4 -free

Thm (Szemerédi - Trotter)

$\forall$ pt/line arrangement w/ $|P| = |L| = n$

$\Rightarrow |I(P, L)| = O(n^{4/3})$

Rmk $ex(n, C_4) = \Theta(n^{3/2})$. Here much better bound for "geometric" graphs

Recall $ex(n, C_6) = \Theta(n^{4/3})$.

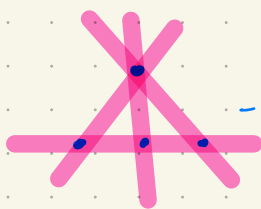
Thm (Solymosi: Dense arr. locally dense)

Given an (n, n) -ux pt/line incidence gr. G ,
if G is C_6 -free $\Rightarrow e(G) = o(n^{4/3})$

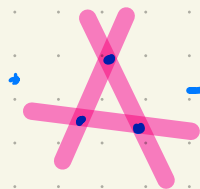
Rephrase Given pt/line arrangement (P, L) with
no 3 pts in general position pairwise lying on some
lines.

$$\Rightarrow |I(P, L)| = o(n^{4/3})$$

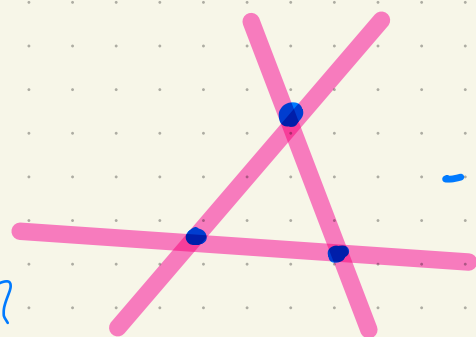
Open • 3-fan



-free $\Rightarrow o(n^{4/3})$?

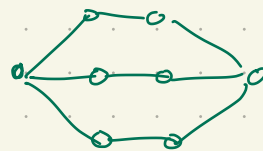


-free $\Rightarrow O(n^{4/3-c})$?



-free

incidence gr
is $\Theta_{3,3}$ -free



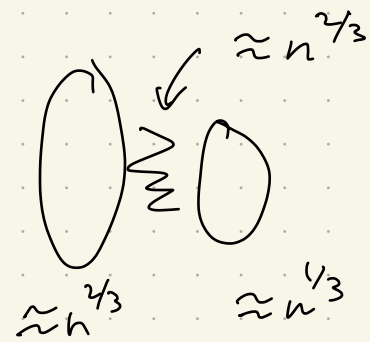
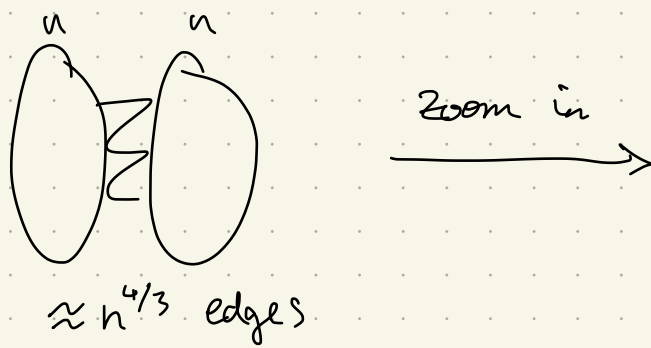
Two ingredients:
- partition lem.
- removal lem.

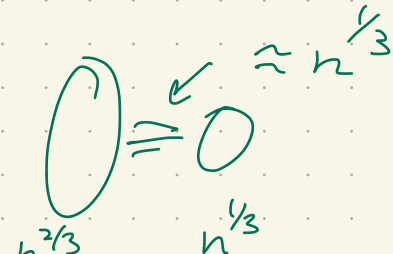
Idea: • graphs obtained from geometric settings
tend to have their edges clustered

$\Rightarrow$ Zoom into exceptionally dense asymm. bip. part.

(Matoušek's partition)

Then use removal lem.



- If we zoom in randomly, $\Rightarrow$ 

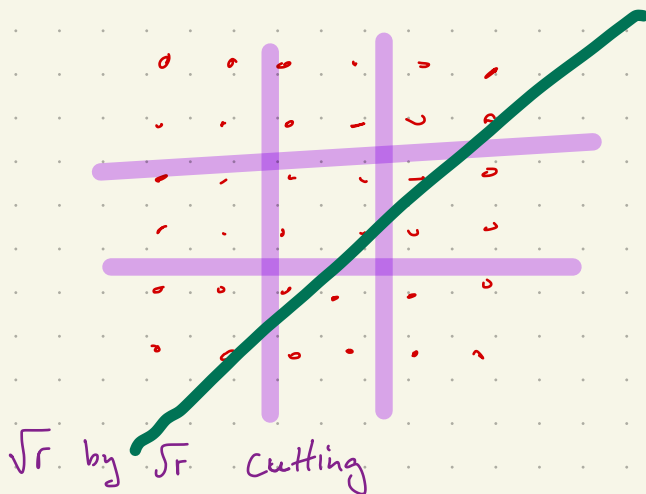
Matoušek partition (d=2)

- Given
- $P \subseteq \mathbb{R}^2$ n pts
 - $2 \leq r < n$
- $\Rightarrow$

$\exists$ partition
 $P = D_1 \cup \dots \cup D_t$

- $\frac{n}{t} \leq |D_i| \leq \frac{2n}{r}$
- any line passes through $O(\sqrt{r})$ sets

Intuitive: imagine pts $\approx$ grid pts



Pf

Suppose the incidence gr G has $\geq c n^{4/3}$ edges.

Take $\beta \ll c$ and set $r = \beta n^{2/3}$.

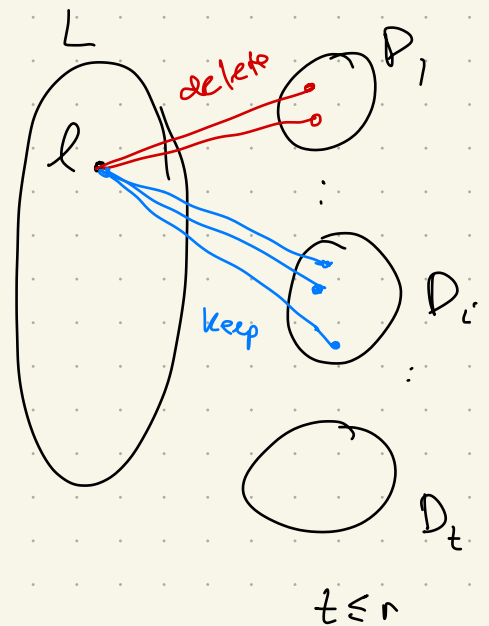
- Apply M. partition $\Rightarrow P = D_1 \cup \dots \cup D_t$

every $v_x \ l \in L$ has neighbors
in $O(\sqrt{r})$ different D_i 's

1) $\forall l \in L \ \forall D_j$

if $\deg(l, D_j) \leq 2$

deletes edges from l to D_j .



$$\# \text{ edges deleted} \leq n \cdot 2 \cdot O(\sqrt{r})$$

$$= O(\beta^{1/2} \cdot n^{4/3}) \leq \frac{1}{2} e(G) \quad \beta \ll c$$

By averaging, $\exists i \in [t]$, s.t.

$$\# \text{ edges from } L \text{ to } D_i \geq \frac{\frac{1}{2} e(G)}{t} \geq \frac{\frac{1}{2} e(G)}{r}$$

$$= \frac{c}{2\beta} n^{2/3}$$

• Look at $[L, D_i]$

$$\frac{n}{r} \leq |D_i| \leq \frac{2n}{r} \approx \frac{1}{\beta} n^{1/3}$$

Build an auxiliary gr. Γ w./

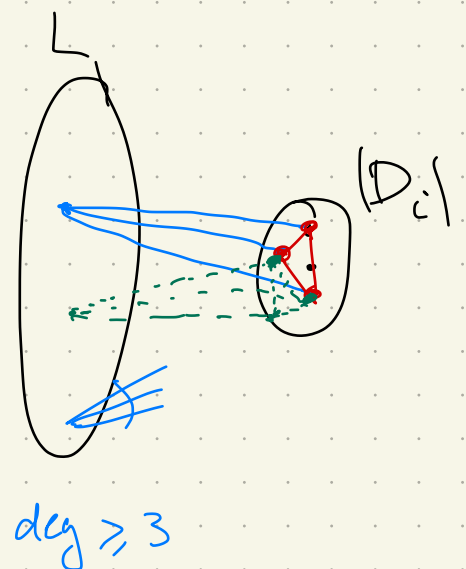
$$V(\Gamma) = D_i$$

$E(\Gamma)$: for each $l \in L$ w./

$$d_l = \deg(l, D_i) \geq 3$$

put $\lfloor \frac{d_l}{3} \rfloor$ vx-disj Δ s in

$$N(l, D_i)$$



- In $\mathcal{T}$, call a triangle L -triangle if it comes from some $l \in L$

all L -triangles are edge-disjoint as

G is C_4 -free

$$\Rightarrow \# \text{ edge-disj } \Delta_s \geq \# L\text{-triangles in } \mathcal{T} = \sum_l \lfloor d_l / 3 \rfloor$$

$$\text{since } d_l \geq 3 \geq \sum_l d_l / 6$$

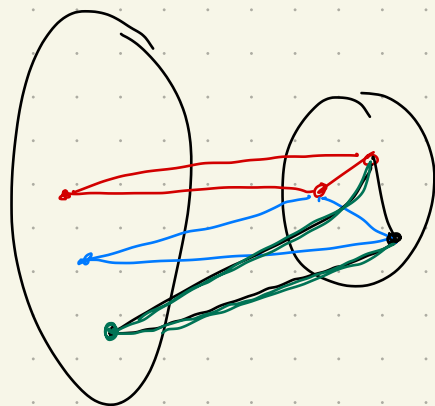
$$\geq \frac{1}{6} \frac{c}{2\beta} n^{2/3} = \Omega(|D_i|^2)$$

• Removal lem $\Rightarrow \mathcal{T}$ has $\Omega(|D_i|^3)$ Δ_s

$$\Omega(|D_i|^3) = \Omega(n) \gg O(n^{2/3}) \geq \# L\text{-triangles}$$

$\Rightarrow \exists$ a non- L -triangle in $\mathcal{T}$.

$\Rightarrow C_6$ in G $\Leftarrow$



Def. $r(G, G) = \min N$ s.t. any 2-edge-col. of K_N
 $\Rightarrow$ monoX. G .

Classical : $2^{n/2} < r(K_n, K_n) < 4^n$

Thm (Chvátal - Rödl - Szemerédi - Trotter)

• $d \in \mathbb{N}$, $\Delta(G) \leq d$

$$\Rightarrow r(G, G) = O_d(|G|)$$

• Multicolor ver. of reg. lem.

For a k -edge-col. gr. G , a p.t.t. $V(G) = V_1 \cup \dots \cup V_r$
is an ε -reg. p.t.t. if

$$\bullet \forall ij \in \binom{[r]}{2}, \quad |V_i| - |V_j| \leq 1;$$

• for all but $\leq \varepsilon \binom{r}{2}$ pairs $ij \in \binom{[r]}{2}$,

(V_i, V_j) is ε -reg. in every color.

Multicol. Reg. Lem. $\forall \varepsilon > 0, k, m \in \mathbb{N}, \exists M = M(\varepsilon, m, k)$

s.t. every k -edge-col. gr. G w/ $n \geq m$ vxs admits an

ε -reg. p.t.t. $V(G) = V_1 \cup \dots \cup V_r$, where $m \leq r \leq M$

Reduced gr. $\leftrightarrow$ reg. p.t.t.

Only difference: assign majority col. to $E(R)$

Thm (Brook's thm) $\forall G \Rightarrow \chi(G) \leq \Delta(G) + 1$.

Pf (Ch-R-Sz-Tr)

Let $n \geq 5r(K_{d+1}, K_{d+1})$, $\varepsilon = 1/n$ and

$$C := \frac{2M}{(\frac{1}{2} - \varepsilon)^d}.$$

Let $N \geq C|G|$ and fix an arbitrary 2-edge-col. of K_N .

Goal: Find a monoX. copy of G .

Apply multicol. reg. lem. to the given 2-edge-col. K_N and let R be the corresp. reduced gr. w/ a 2-edge-col. indicating the majority color.

• As $\leq \varepsilon \binom{n}{2}$ irreg pairs, R is almost complete:

$$e(R) \geq (1 - \varepsilon) \binom{n}{2} > \left(1 - \frac{1}{m/3 - 1}\right) \frac{r^2}{2}$$

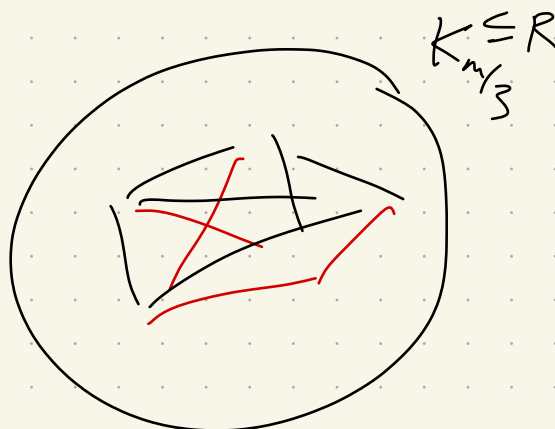
• Turán's thm $\Rightarrow R$ contains a clique $K_{m/3}$

• $m/3 > r(K_{d+1}, K_{d+1})$

$\Rightarrow \exists \text{ mono } \chi, K_{d+1} \subseteq K_{m/3} \subseteq R$

• K_{d+1} is a hom. image of G

as $\Delta(G) \leq d$ by Brook's thm



• Embedding Lem: $\text{mono } \chi K_{d+1} \subseteq R \rightsquigarrow \text{mono } \chi G$
in K_N .

