

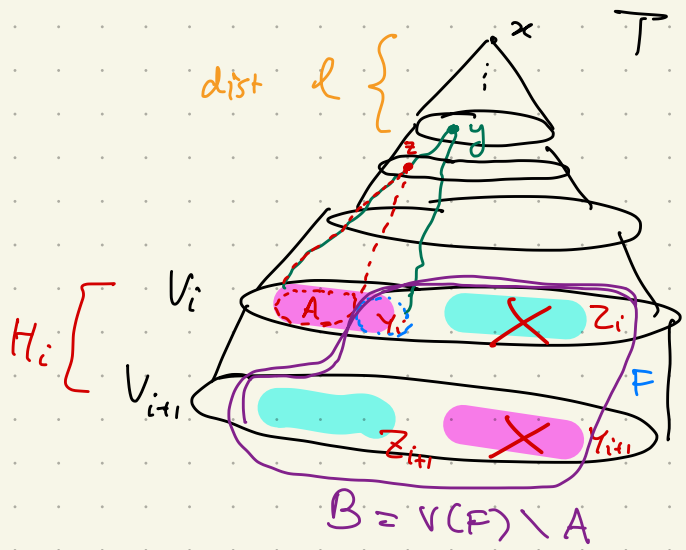
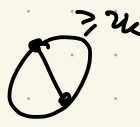


# Lecture 15

• More details on missing cases

$\Theta_k$ -graph $_1^F$  w/ bipartition  $Y \cup Z$

•  $\Theta_k$ -graph $_1^F$  is bipartite  
and  $\delta(F) \geq 2$ , connected!



If, wlog,  $Y \cap V_i \neq \emptyset$

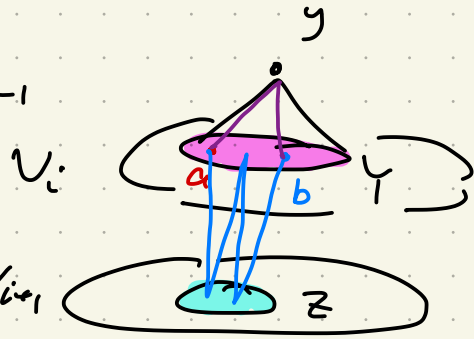
$\Rightarrow$

we must have  $Y \subseteq V_i, Z \subseteq V_{i+1}$

• Define  $y$  as before. Case 1 ;  $y \notin V_{i-1}$ , define  $z$  as before, ... do the same pf.

Case 2 (Missing case) if  $y \in V_{i-1}$

then take an arbitrary  $a \in Y$  and let



$A = \{a\}$ , and  $B = F \setminus A$

Again  $\delta(F) \geq 2 \Rightarrow (A, B)$  is NOT a bip. of  $F$

(as  $B$  contains edges)

$\Rightarrow \exists A, B$ -paths $_1^P$  of length  $2k-2$  (even)

$P$  has endpoints  $a$  and  $b$  both in  $Y$ .

$P + ay + by \Rightarrow C_{2k} \quad \square$



## § Regularisation.

A lem of Erdős and Simonovits allows us to work with almost regular graphs for bipartite Turán problem.

We say a graph  $G$  is  $K$ -almost regular if

$$\Delta(G) \leq K \cdot \delta(G)$$

Lem (E-Sim) Let  $0 < \epsilon < 1$ ,  $c > 0$  and  $n$  be suff.

large. Let  $G$  be an  $n$ -vertex graph with

$e(G) \geq c \cdot n^{1+\epsilon} \Rightarrow$  Then  $G$  contains an  $m$ -vertex

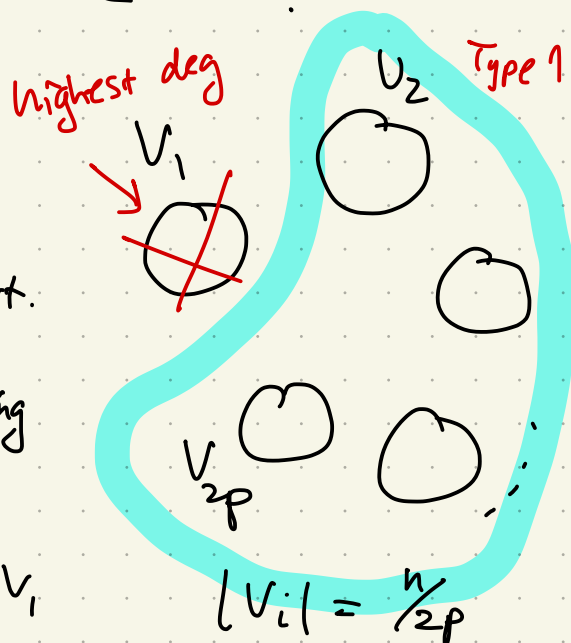
$K$ -almost regular subgraph  $G'$  w/  $m \geq n^{\frac{\epsilon - \epsilon^2}{4 + 4\epsilon}}$  and

$e(G') \geq \frac{2c}{5} m^{1+\epsilon}$ , where  $K = 20 \cdot 2^{\frac{1}{\epsilon^2} + 1}$ .

Pf (Sketch, details = exercise)

Let  $p = 2^{\frac{1}{\epsilon^2} + 1} = \frac{K}{20}$  and take a  $2p$ -equipart.

$V_1 \cup \dots \cup V_{2p}$  of  $V(G)$  with  $V_1$  containing the highest deg. vxs.



- $G$  type 1: if  $\leq$  half edges incident to  $V_1$

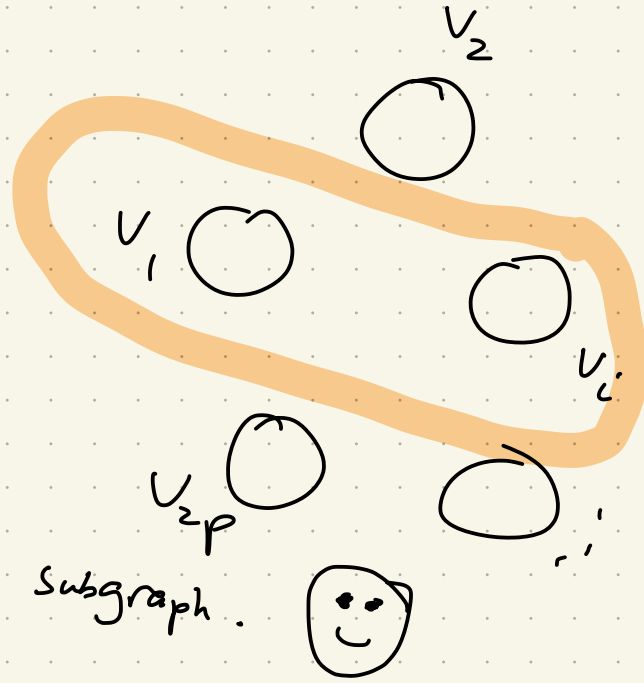
$\uparrow$  - delete  $V_1$

$\uparrow$  - repeated delete vxs of  $\deg < \frac{c}{10} n^\epsilon$  until no such vxs

$\exists$  exist  $\Rightarrow$  desired  $G'$ . 😊

•  $G$  type 2: if  $\geq$  half edges incident to  $V_i$  Type 2

- Pigeonhole  $\Rightarrow \exists v_i$  s.t.  
 $e_G(v_i, v_i) \geq \frac{1}{2p} e(G)$   
 - iterate the same analysis  
 on  $G[v_i, v_i]$



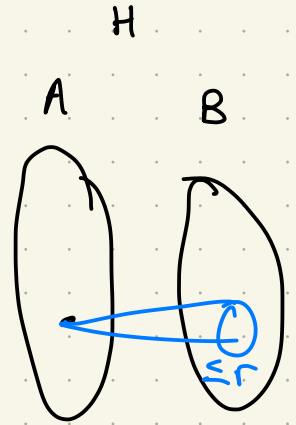
It can be shown that this process terminates at a large type 1 subgraph. 😊

§ Bip. graphs w/ bdd deg.

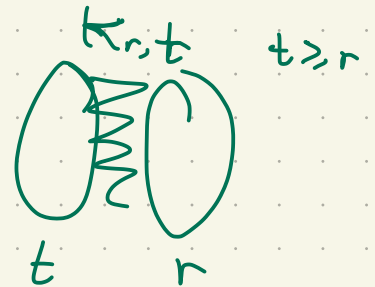
Thm (Füredi, Alon-Krivelevich-Sudakov)

Let  $H$  be a bip. graph w/ bipartition  $A \cup B$  where each  $v_x$  in  $A$  has  $\deg \leq r$  in  $B$ . Then  $\exists c = c(H)$  s.t.

$$ex(n, H) \leq c \cdot n^{2 - \frac{1}{r}}$$



Remark This thm generalise K-S-T.



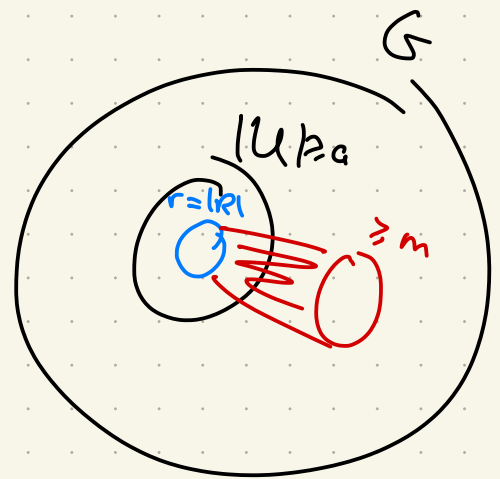
# § 1st pf via dependent random choice lem (DRC)

The following basic lem finds in a dense graph a large set  $U$  s.t. all small subsets of  $U$  have many common neighbors.

Lem Let  $a, d, m, n, r \in \mathbb{N}$ . Let  $G = (V, E)$  be an  $n$ -vertex graph with average deg  $d$ . If  $\exists t \in \mathbb{N}$  s.t.

$$\frac{d^t}{n^{t-1}} - \binom{n}{r} \left(\frac{m}{n}\right)^t \geq a$$

$\Rightarrow$  then  $G$  contains a subset  $U$  of size  $\geq a$  s.t. every  $r$ -subset  $R \subseteq U$  has  $\geq m$  common neighbors.



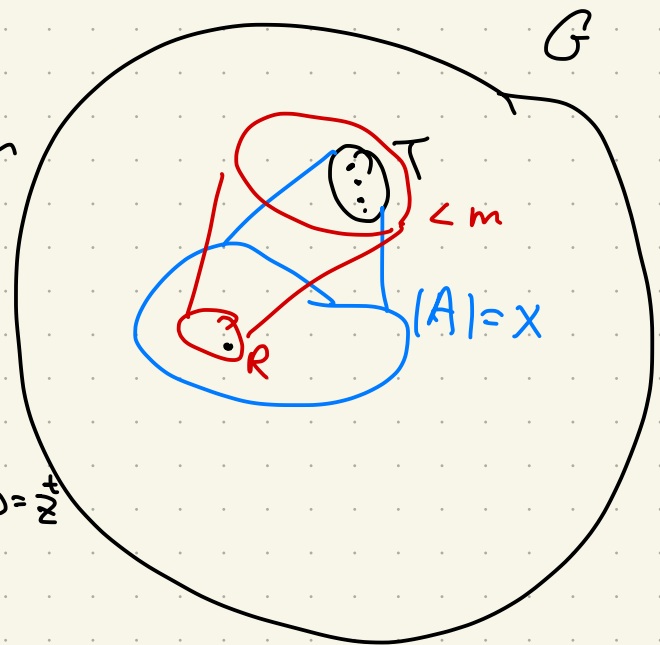
PF: • Pick a set  $T$  of  $t$  vs of  $V$  uniformly at random w./ repetition

Let  $A = N^*(T) = \bigcap_{t \in T} N(t)$  and let

$$X = |A|$$

• Linearity of  $\mathbb{E}$

$$\Rightarrow \mathbb{E} X = \sum_{v \in V} \left(\frac{d(v)}{n}\right)^t \stackrel{\text{convexity of } f(z) = \frac{z^t}{2}}{\geq} \frac{d^t}{n^{t-1}}$$



Let  $Y = \#$   $r$ -sets  $R \subseteq A$  w./  $|N^*(R)| < m$ .

For a set  $\bigwedge^R$  w./  $< m$  common neighbors,

$$\mathbb{P}(R \subseteq A) = \mathbb{P}(T \in N^*(R)) = \left(\frac{|N^*(R)|}{n}\right)^t < \left(\frac{m}{n}\right)^t$$

$$\Rightarrow \mathbb{E}Y \leq \binom{n}{r} \cdot \left(\frac{m}{n}\right)^t$$

$$\Rightarrow \mathbb{E}(X - Y) \geq \frac{d^t}{n^{t-1}} - \binom{n}{r} \left(\frac{m}{n}\right)^t \geq a$$

Therefore,  $\exists$  a choice of  $\tau$  s.t.  $X - Y \geq a$

Delete one  $u_x$  from each 'bad'  $R$  in  $A$

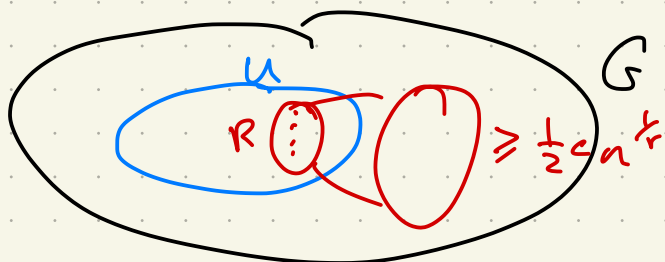
$\Rightarrow |U| \geq X - Y \geq a$  w./ the desired property. 😊

Cor (Exercise) Let  $G$  be an  $n$ - $u_x$  graph w./

$\geq cn^2$  edges and  $r \in \mathbb{N}$ .

$\Rightarrow \exists$  a subset of  $u_x$ s  $U$  of size  $|U| \geq \frac{1}{2} c^r n$

s.t. all  $r$ -sets in  $U$  have  $\geq \frac{1}{2} c n^{1/r}$  common neighbors.



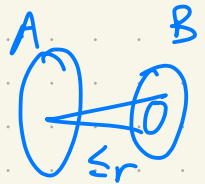
Exercise 1 Use the above Cor. to prove

Füredi - Alon - Krivelevich - Sudakov.

• Use dependent random choice lem to prove the following.

Exercise  $G \begin{cases} n\text{-vx} \\ cn^2 \text{ edges} \end{cases} \Rightarrow$  1-subdivision of  $K_t$ ,  
 where  $t \geq c^{3/2} n^{1/2}$ .

Rmk Fox-Sudakov survey on DRC.



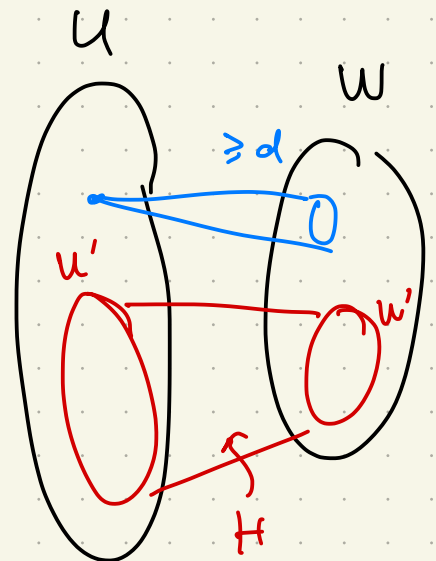
§ 2<sup>nd</sup> pf of F-AK-S via random zooming

Let  $H$  be <sup>an  $h$ -vertex</sup> bip. graph w/ bip.  $A \cup B$  s.t. every  $vx$  in  $A$  has  $\deg \leq r$  in  $B$ .

Thm 1 Let  $d \geq 40$  and  $G$  be a bip. graph w/ bip.  $U \cup W$  s.t. every vertex in  $U$  has  $\deg d$  in  $W$ .

If  $\frac{1}{2|W|} \cdot \left(\frac{|U|}{4h}\right)^{1/r} \cdot \frac{d}{2} \geq 20$

$\Rightarrow$  then  $H \subseteq G$



PF (Thm 1  $\Rightarrow$  F-A-K-S:  $ex(n, H) = O(n^{2-\frac{1}{r}})$ )

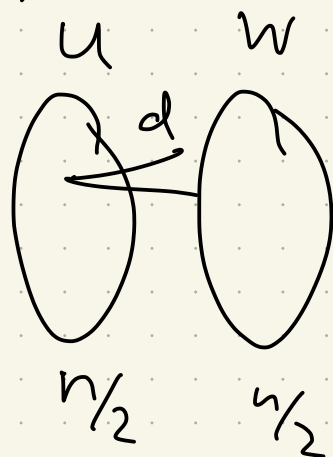
$$d = c' n^{1-\frac{1}{r}}$$

$$\frac{1}{2|W|} \cdot \left(\frac{|U|}{4h}\right)^{\frac{1}{r}} \cdot \frac{d}{2}$$

$$= \frac{1}{2 \cdot \frac{n}{2}} \cdot \left(\frac{n}{8h}\right)^{\frac{1}{r}} \cdot \frac{c'}{2} n^{1-\frac{1}{r}} \geq 20$$

as long as  $c' \gg h, r$

Exercise: Prove F-A-K-S using Thm 1.



PF Claim  $\exists U' \subseteq U, W' \subseteq W$  s.t.

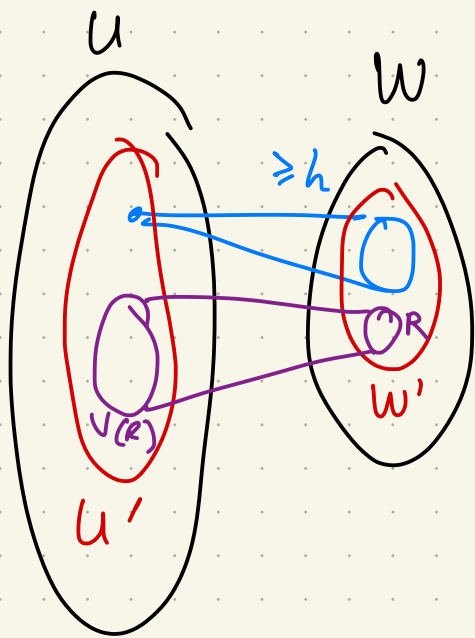
- $|U'| \geq h \cdot |W'|^r$  (very skewed)
- every  $u_x$  in  $U'$  has  $\deg \geq h$  in  $W'$

We shall embed  $H$  in  $G[U', W']$

- Take a maximal collection

$\mathcal{I} \subseteq \binom{W'}{r}$  for which there is

$V(R) \subseteq U'$  for every  $R \in \mathcal{I}$  s.t.



$V(R)$  is completely joined to  $R$  and  $|V(R)| = h$ .

• As  $|U'| \geq h \cdot |W'|^r > h \cdot \binom{|W'|}{r}$

$\Rightarrow \exists$  a vertex  $b \in U'$  s.t.  $b \notin V(R)$  for any  $R \in \mathcal{I}$ .

Take  $B' \subseteq N(b, W')$ ,  $|B'| = h$

Embed the B side of  $H$  to  $B'$ .

greedily embed  $A$  using

$V(R)$ ,  $R \in \mathcal{I}$ .

