

Lecture 15

Announcement: No class next week.

Weak ϵ -net

Fish \Leftrightarrow polytope
convex body

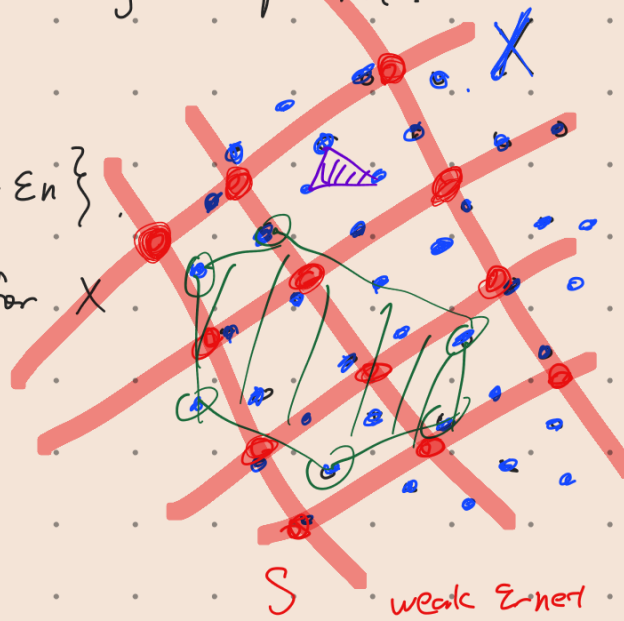
Def: Let $X \subseteq \mathbb{R}^d$ be a set of n pts in general position.

For $\epsilon > 0$, define

$$\mathcal{F}_\epsilon = \{ \text{conv } Y : Y \subseteq X, |Y| \geq \epsilon n \}$$

A finite set $S \subseteq \mathbb{R}^d$ is a **weak ϵ -net** for X

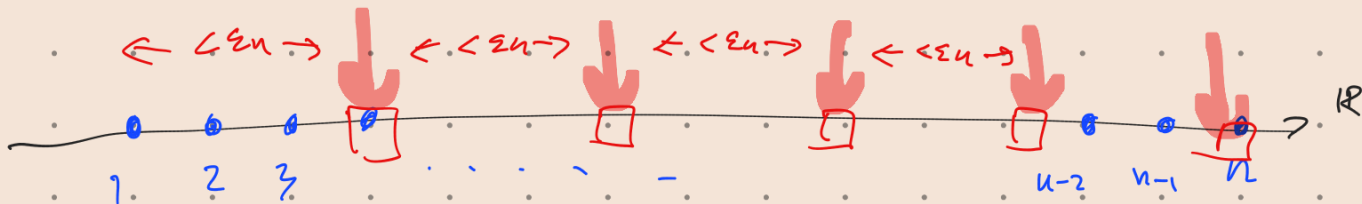
if $S \cap F \neq \emptyset \quad \forall F \in \mathcal{F}_\epsilon$.



Rank • Note that pts in the weak ϵ -net are not necessarily from X

• Informally, ^{weak} ϵ -net only 'captures' the "chubby fish".

Ex \mathbb{R}^1 , $X = \{1, 2, 3, \dots, n\}$



• Here in this example, the chubby fish are the intervals containing $\geq \epsilon n$ pts in X .

• Take $\{ \lfloor \epsilon n \rfloor, 2 \lfloor \epsilon n \rfloor, \dots, t \lfloor \epsilon n \rfloor \}$ where $t = \lfloor \frac{1}{\epsilon} \rfloor$

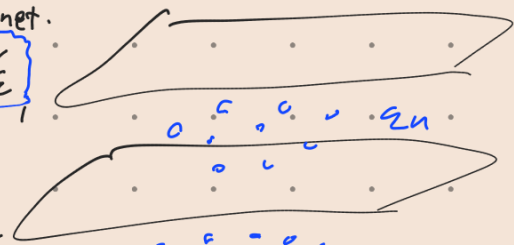
\Rightarrow weak ϵ -net of size $\leq \frac{1}{\epsilon}$.

Rank: Cannot do better. In \mathbb{R}^d ,

weak ϵ -net.
 $\geq \frac{1}{\epsilon}$

Take $\frac{1}{\epsilon} + 1$ parallel hyperplanes in \mathbb{R}^d

and place ϵn pts between two consecutive



ones $\Rightarrow \geq \frac{1}{\epsilon}$ pts in any

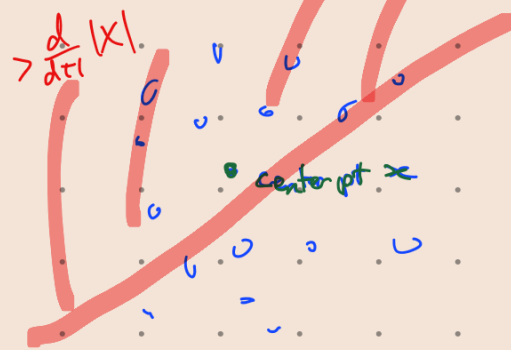
weak ϵ -net. as we need ≥ 1 pts between any two consecutive hyperplane.

Thm: (weak ϵ -net thm) \forall finite $X \subseteq \mathbb{R}^d$ and $\forall \alpha \times \epsilon < 1$,

$\Rightarrow \exists$ a weak ϵ -net of size $\leq c_d \frac{1}{\epsilon^{d+1}}$, where $c_d = 2^{d+1} (d+1)^d$

Rmk: - Connection to Helly: Center pt

$\forall \epsilon > \frac{d}{d+1}$, Helly thm \Rightarrow there is a weak ϵ -net of size 1.



• Best lower bound [Bukh - Matoušek - Nivasch] $\geq c_d \frac{1}{\epsilon} (\log \frac{1}{\epsilon})^{d-1}$

Upper bound • Matoušek - Wagner $\leq \tilde{O}(\frac{1}{\epsilon^d})$

• Rubinfeld $\leq O(\epsilon^{-d-1/2-\epsilon})$

Idea:

• each **chubby fish** contains many X -simplices $(\geq \binom{\epsilon n}{d+1})$

By **selection lem**

• deepest pt inside pieces many X -simplices $(\geq c_d' \binom{\epsilon n}{d+1})$

• so $\frac{\binom{n}{d+1}}{c_d' \binom{\epsilon n}{d+1}} \leq c_d \frac{1}{\epsilon^{d+1}}$ rounds, we pierce all X -simplices

pf: • May assume $\epsilon \leq \frac{1}{2(d+1)}$ (since $\forall \epsilon' > \epsilon$, any weak ϵ -net is also a weak ϵ' -net)

• May assume $n \geq \frac{2(d+1)}{\epsilon}$ (o.w. $c_d \frac{1}{\epsilon^{d+1}} \geq n$, we can take the whole X to be a weak ϵ -net)

Construct a weak ϵ -net iteratively.

Initiate $S_0 = \emptyset$ and $f_0 = \binom{X}{d+1}$

if $F \in \mathcal{F}_\varepsilon = \{ \text{conv } Y : Y \subseteq X, |Y| \geq \varepsilon_n \}$

s.t. $F \cap S_{i-1} = \emptyset$

then apply selection lemma to get a pt x_i

packing $\geq \frac{1}{(d+1)^d} \binom{|I|}{d+1}$ many X -simplices in F

Set $S_i = S_{i-1} \cup \{x_i\}$

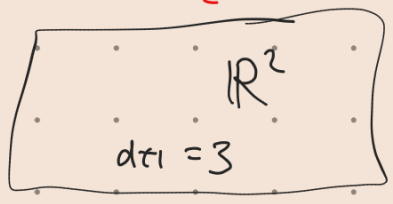
else terminate.

By construct, the set S_t at the end is a weak ε -net.

and $t = |S_t| \leq \frac{\binom{n}{d+1}}{\frac{1}{(d+1)^d} \binom{\varepsilon n}{d+1}} \leq C_d \frac{1}{\varepsilon^{d+1}}$

$\varepsilon \leq \frac{1}{2(d+1)}$
 $n \geq \frac{2(d+1)}{\varepsilon}$

Lem: Let C_1, \dots, C_{d+1} be convex sets in \mathbb{R}^d , if \forall nontrivial partition $[d+1] = A \cup B$, $\bigcup_{i \in A} C_i$ and $\bigcup_{j \in B} C_j$ can be strictly separated by a hyperplane



\Rightarrow then there is no hyperplane simultaneously intersecting C_1, \dots, C_{d+1}

Moreover, if $C_i = \text{conv } X_i$ for some set $X_i \subseteq \mathbb{R}^d$ (finite)

$\Rightarrow (X_1, \dots, X_{d+1})$ has the same type transversal.

Idea (for some type lem) is to repeatedly apply Helly-sandwich to obtain large subsets that are well-separated.

For each partition $A \cup B = [d+1]$ we also a function $f: C \rightarrow \mathbb{R}$

from a hyperplane (from Haim-Sandwich).

We need to apply it for $\leq 2^{d+1}$ times. $\Rightarrow |Y_i| \geq 2^{-2^{d+1}} |X_i|$

